NAVAL POSTGRADUATE SCHOOL
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THESIS

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Implementation of Multi-Frequency Modulation
with Trellis Encoding and Viterbi Decoding
Using a Digital Signal Processing Board

by

John W. Wisniewski
June 1991

Thesis Advisor: P.H. Moose

92-03678

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92 2 12 159

REPORT DOCUMENTATION PAGE

Form Approved
OMB No. 0704-0188

1a. REPORT SECURITY CLASSIFICATION UNCLASSIFIED	1b. RESTRICTIVE MARKINGS
2a. SECURITY CLASSIFICATION AUTHORITY	3. DISTRIBUTION / AVAILABILITY OF REPORT
2b. DECLASSIFICATION / DOWNGRADING SCHEDULE	Approved for public release; distribution is unlimited

4. PERFORMING ORGANIZATION REPORT NUMBER(S)	5. MONITORING ORGANIZATION REPORT NUMBER(S)

6a. NAME OF PERFORMING ORGANIZATION Naval Postgraduate School	6b. OFFICE SYMBOL (If applicable) 32	7a. NAME OF MONITORING ORGANIZATION Naval Postgraduate School
6c. ADDRESS (City, State, and ZIP Code) Monterey, California 93943-5000		7b. ADDRESS (City, State, and ZIP Code) Monterey, California 93943-5000

8a. NAME OF FUNDING / SPONSORING ORGANIZATION	8b. OFFICE SYMBOL (If applicable)	9. PROCUREMENT INSTRUMENT IDENTIFICATION NUMBER
8c. ADDRESS (City, State, and ZIP Code)		10. SOURCE OF FUNDING NUMBERS

PROGRAM ELEMENT NO.	PROJECT NO.	TASK NO.	WORK UNIT ACCESSION NO.

11. TITLE (Include Security Classification)
IMPLEMENTATION OF MULTI-FREQUENCY MODULATION WITH TRELLIS ENCODING AND VITERBI DECODING USING A DIGITAL SIGNAL PROCESSING BOARD

12. PERSONAL AUTHOR(S)
WISNIEWSKI, John William

13a. TYPE OF REPORT Master's Thesis	13b. TIME COVERED FROM _____ TO _____	14. DATE OF REPORT (Year, Month, Day) 1991 June	15. PAGE COUNT 95

16. SUPPLEMENTARY NOTATION The views expressed in this thesis are those of the author and do not reflect the official policy or position of the Department of Defense or the U.S. Government.

17. COSATI CODES			18. SUBJECT TERMS (Continue on reverse if necessary and identify by block number)
FIELD	GROUP	SUB-GROUP	Communication; Multi-Frequency Modulation; Trellis coding; Viterbi decoding

19. ABSTRACT (Continue on reverse if necessary and identify by block number)

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20. DISTRIBUTION / AVAILABILITY OF ABSTRACT ☒ UNCLASSIFIED/UNLIMITED ☐ SAME AS RPT ☐ DTIC USERS	21. ABSTRACT SECURITY CLASSIFICATION UNCLASSIFIED	
22a. NAME OF RESPONSIBLE INDIVIDUAL MOOSE, P. H.	22b. TELEPHONE (Include Area Code) (408)-646-2838	22c. OFFICE SYMBOL EC/Me

DD Form 1473, JUN 86   Previous editions are obsolete.   SECURITY CLASSIFICATION OF THIS PAGE

S/N 0102-LF-014-6603

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by

John W. Wisniewski
Lieutenant, United States Navy
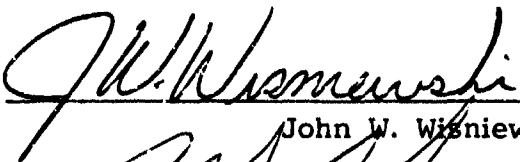B.S.E.E., Virginia Military Institute

Submitted in partial fulfillment
of the requirements for the degree of

MASTER OF SCIENCE IN ELECTRICAL ENGINEERING

from the

NAVAL POSTGRADUATE SCHOOL
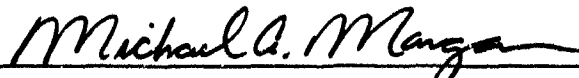June 1991

Author: _____
John W. Wisniewski

Approved by: _____
P. H. Moose, Thesis Advisor

_____
T. T. Ha, Second Reader

_____
Michael A. Morgan, Chairman
Department of Electrical and Computer Engineering

ABSTRACT

Multi-Frequency Modulation has been the topic of several papers at NPS. In past systems the majority of time required for the generation of the MFM signal was due to the software routine used to implement the FFT. In this report a Digital Signal Processor was used to reduce the time needed to generate the FFT. The use of Trellis coding and Viterbi decoding on a Digital Signal Processor was also investigated. Assembly language programs for three encoder/decoder systems were developed. The first uses a 16 QAM signal, the second uses a 2/3 rate convolutional encoder and Viterbi decoder and the third uses the V.32 convolutional encoder and a Viterbi decoder.

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I. INTRODUCTION

There has long been an interest in frequency division multiplexing as a means of combatting impulsive noise, avoiding equalization and making fuller use of the bandwidth available. The principle for Multi-Frequency Modulation (MFM) was first used over thirty years ago in the Collins Kineplex system. Since that time MFM has been used under many names:

- multiplexed Quadrature Amplitude Modulation (QAM)

- orthogonal Frequency Division Multiplexing (FDM)

- dynamically assigned QAM

- multicarrier modulation

A parallel or multiplexed data system offers the potential to reduce some of the problems of serial systems. In order to increase the data rate in a serial system you must utilize higher order modulation or decrease the symbol interval, increasing the bandwidth, at the risk of degradating the performance of the system. In a parallel data system several sequential streams of data are transmitted simultaneously. In a classical parallel data system there are N non-overlapping subchannels with each data element occupying only a small portion of the available bandwidth.

With the continuing development of Digital Signal Processing technology there has been renewed interest in MFM for possible use in:

- General Switched Telephone Network (GSTN)

- 60-108 Khz Frequency-division Multiplexed (FDM) group band

- Cellular radio

- High speed data for transmission on the high-rate digital subscriber line (HDSL)

Multi-Frequency Modulation (MFM) can easily be implemented on a computer due to the minimal requirements of hardware to construct a transmitter or receiver. The primary components of a MFM transmitter/receiver are D/A, A/D converters and a method for computing an FFT. MFM has been the topic of several papers and projects at NPS using various techniques for data acquisition and modulation. One of the primary problems with implementing an MFM system in past projects, has been the time required to perform an FFT using current software routines. With the continuing development of Digital Signal Processing (DSP) chips there exist a great number of DSP boards complete with D/A, A/D converters, filters, and memory available for industry standard computers. This provides the capability to purchase an off the shelf item that has all the necessary components to implement an MFM transmitter/receiver. Utilizing a DSP board for the computation of the FFT's required for MFM provides a tremendous performance gain over the software routines that were used in previous projects. Once the programs are downloaded, the performance of the DSP board is independent of the performance of the host computer.

General purpose Digital Signal Processor chips, such as the Motorola DSP 56001 (which was used in the development of the programs included), have an architecture ideally suited to the rapid calculation of FFTs. The 56001 is a fixed point Digital Signal Processor which has three internal memory and address busses which allow the execution of instructions with

parallel data moves to the X and Y memory locations. This allows the rapid execution of complex computations. There are also 256 locations of X and Y Rom that contain Mu-law and A-law expansion tables as well as a full four quadrant sine wave table.

The first objective of this project is to investigate the use of a DSP for implementing an MFM system and reduce the time required to generate a MFM signal. The second objective was to investigate the use of trellis coding and Viterbi decoding using a Digital Signal Processor.

II. THEORY

A. MULTI-FREQUENCY MODULATION

A Multi-Frequency Modulated signal has a packet structure and is comprised of time and frequency slots.(Figure 1)
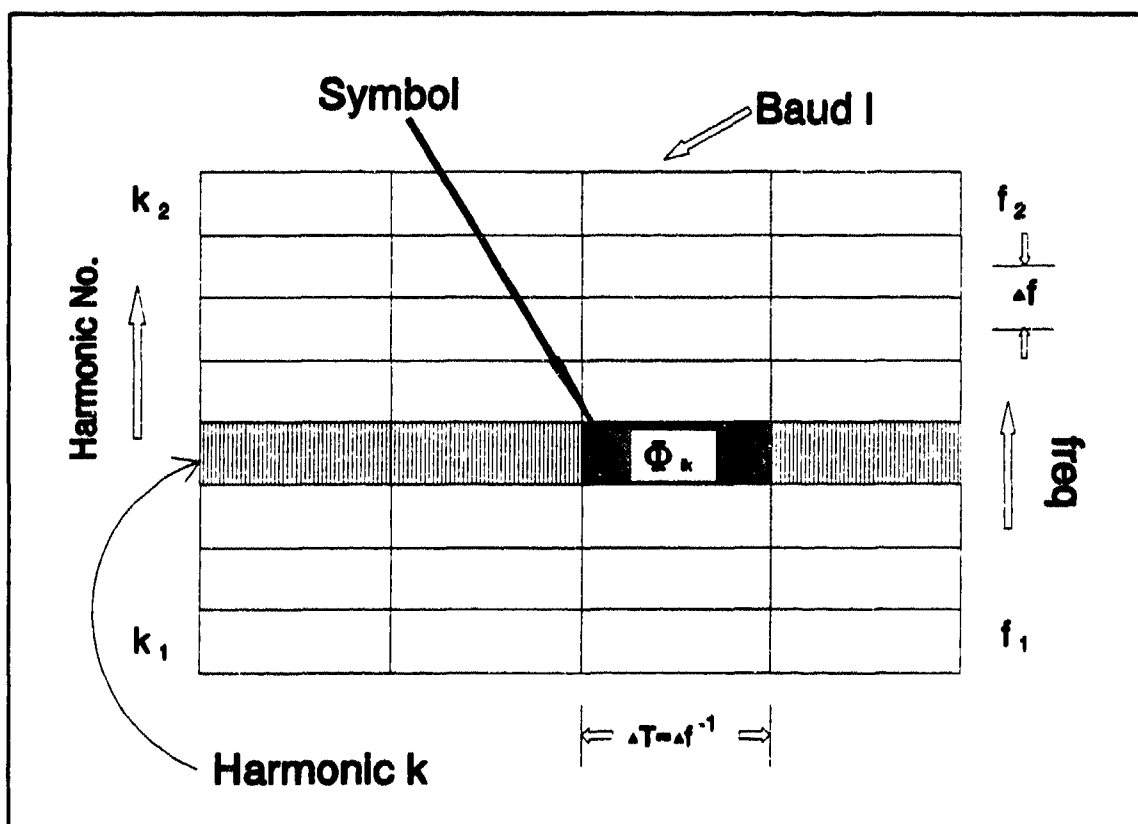
Figure 1 MFM Signal Packet (from Ref. 1: p.3)

The following definitions will be used for discussing MFM: [Ref. 1:pp 5-6]

- K: Number of MFM tones

- T: Packet length in seconds

- ΔT: Baud length in seconds

- k_x: Baud length in number of samples

- L: Number of bauds per packet

- Δt: Time between samples in seconds

- $f_x = 1/\Delta t$: Sampling or clock frequency for D/A and A/D conversion in Hz

- $\Delta f = 1/\Delta T$: Frequency spacing (minimum) between MFM tones

- Φ_{lk}: Symbol set. Phase of the kth tone in the lth baud

- A_{lk}: Amplitude of the kth tone in the lth baud

Each MFM packet consists of L bauds of K tones. The packet construction for the lth baud is given by:[Ref 1: pp 6-7]

$$x_l(u) = \sum_k x_{lk}(u) \qquad (1)$$

where, the analog representation of each tone during the lth baud is given by:

$$x_{lk}(u) = A_{lk}\cos(2\pi k\Delta fu + \phi_{lk}) \qquad ; 0 \leq u \leq \Delta T \qquad (2)$$

The time ,u , is referenced from the beginning of the baud. The time at the beginning of the baud is defined as t_0 and the time at any given point in the packet is $t = t_0 + l\Delta T + u$.

By sampling (1) and (2) at intervals $\Delta t = 1/f_x$ a discrete time sampled version for the lth baud can be found:

$$x_l(n) = \sum_k x_{lk}(n) \qquad (3)$$

where the discrete time version of each tone during the lth baud is given by:

$$x_{lk}(n) = A_{lk}\cos(2\pi kn/k_x + \phi_{lk}) \; ; \; 0 \leq n \leq k_x - 1 \qquad (4)$$

The value, n, is the discrete time referenced to the beginning of the baud. From the Sampling Theorem the maximum frequency must be less than $f_x/2$. There are a maximum of $k_x/2$ tones available spaced at intervals of Δf between dc and $f_x/2 - \Delta f$. The value of k is in the range from 0 and $k_x/2$, where k is the harmonic number.

By taking the Discrete Fourier Transform (DFT) of (3):

$$X_l(k) = \sum_k X_{lk}(k) \qquad (5)$$

The discrete time signal (3) can be generated by the k_x point IDFT of (5):

$$x_l(n) = IDFT[X_l(k)] \qquad (6)$$

It is easily seen that the lth baud is generated by taking the IDFT of a complex valued array of length k_x. The values in the array, $x_l(n)$ are the discrete time samples of the analog transmit signal. The generation of the lth baud is completed by clocking the k_x samples out at f_x samples/sec. The entire packet is completed by an L fold repetition of the procedure.

B. TRELLIS CODED MODULATION

Trellis coded modulation evolved as a combination of coding and modulation techniques for digital transmission over band limited channels. The primary advantage for using trellis coded modulation over other coding schemes is that significant coding gains can be achieved without compromising bandwidth efficiency [Ref 2:pp 5-12]. Trellis coded

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modulation schemes employ redundant nonbinary modulation in combination with a finite-state encoder to govern the selection of modulation signals and to generate coded signal sequences. A simple four state scheme can improve the reception of a digital transmission by as much 3db in additive white gaussian noise. If more complex coding schemes are used gains of 6db or more may be achieved. In order to achieve the potential gains of a trellis coded signal a soft decision decoder, such as a Viterbi decoder, must be utilized in the receiver. These gains can also be achieved without reducing the information rate or increasing the bandwidth as required by conventional error correction schemes.

For low to medium data rates (≤ 4800 bits/s), signaling methods that use independent symbol-by-symbol transmission are adequate to provide acceptably low error rates over voice grade circuits. When data rates increase in speed (≥ 9600 bits/s) the same is not true, QAM and optimal signal sets fail to provide acceptably low error rates.[Ref 3:pp 648-649]

There are many factors which comprise linear and nonlinear distortion on a voiceband channel including phase jitter, frequency offset and additive noise. It was shown in that, relative to the uncoded system, trellis codes with four and eight states provide marked improvement in performance with respect to additive noise, second and third harmonic distortion, phase jitter, impulse noise and other channel imparements.[Ref 3:pp 649-651]

For an uncoded system the binary data transmit rate is equal to m/T bits/sec, where m is the number of bits/symbol and $1/T$ the symbol rate. In a conventional QAM system (uncoded) there are 2^m discrete symbol points (amplitude and/or phase levels) with each successive symbol transmitted

independently. The error performance depends on the minimum distance between the signal points (the larger the distance, the lower the error rate). The minimum distance at the transmitter is limited by the average power allowed on the circuit and the choice of the signal points. Even with the use of in-phase and quadrature components the minimum distance (d_{min}) between points decreases with the increase in the number of symbols, m, and constant average power. This results in a degradation at higher transmission rates assuming that there is a constant symbol rate.

The objective of coding is to increase the effective minimum distance between the signals without increasing the average power. One method of accomplishing this is with the use of convolutional encoder. In a convolutional encoder there are m current bits, and v past bits used to develop the codeword. The v past bits define the state of the encoder and are operated on to produce m+j bits, the rate of the code is described as m/m+j. The m+j bits require 2^{m+j} discrete channel symbols. Using a convolutional encoder the minimum distance between symbols is no longer the measure of error performance, performance is now a measure of the minimum distance between the allowed transition of symbols from one state to another.[Ref 3:p 649] A convolutional encoder using a shift register to provide two past bits for a 2/3 rate convolutional code is shown in Figure 2. The allowed transition between states is shown in Figure 3. A convolutional code can be described as an (n,k,m) code, where n is the number of encoded bits, k the number of information bits and m the number of past bits used for encoding.
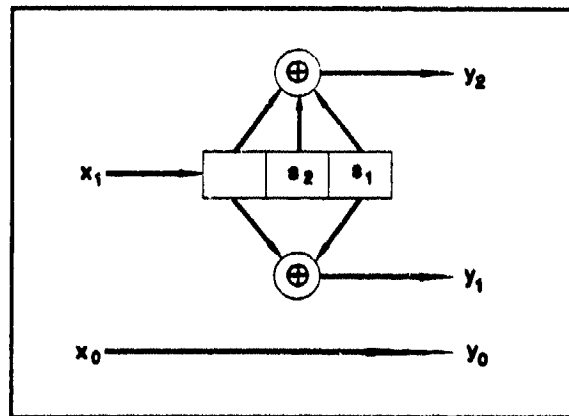
Figure 2 2/3 Rate Convolutional Encoder (from Ref 3:pp 651)

The input and output relationship are depicted by the branches on the trellis diagram. The upper branch depicts the transition with x_0 set to "0", the lower branch depicts the transition with x_0 set to "1". Each node of the trellis diagram represents one of four states created by the past bits s_1 and s_2.
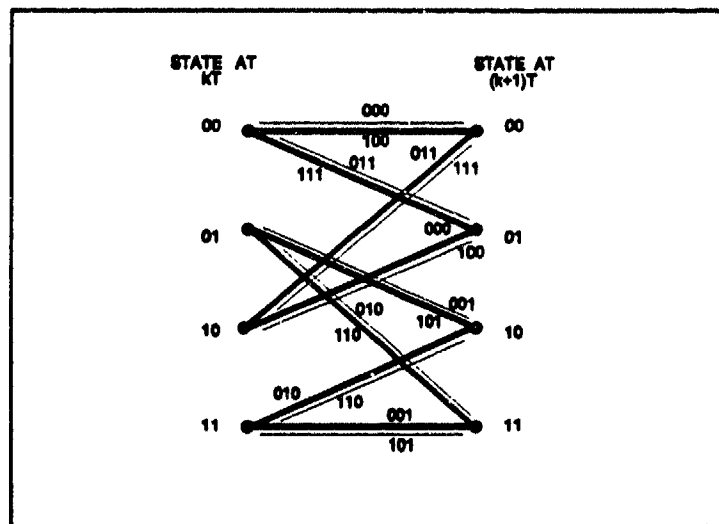
Figure 3 Trellis for 2/3 rate convolutional code.(from Ref 3:pp 651)

In order to achieve the optimum decoding gains from the use of convolutional encoders the decoder must use a trace back routine to find

the most probable path through the trellis. The rules for bit to symbol mapping for coded systems:[Ref 3:pp 650]

1. All parallel transitions in the trellis structure receive maximum possible Euclidean distance in the signal constellation.

2. All transitions diverging from a merging into a trellis state receive maximum possible Euclidean distance.

C.   VITERBI DECODER

The Viterbi decoder utilizes a maximum likelihood sequence estimation method to decode the incoming data stream. A predetermined measure is used to determine the symbol sequence which is closest to the received symbol sequence. At any time, k, the shortest path, called the survivor, entering each state (node) of the trellis is retained. To proceed to time k+1, all time k survivors are extended by computing the metrics (lengths) of the extended path segments based on the calculated branch metrics, dependent on the branch symbols in the trellis and the value of the received sample.[Ref 3:pp 648-649] An example of a trellis and a path through the trellis is shown in Figure 4. The metrics of the remaining paths are computed and the shortest path retained. The shortest length into each state, the (k+1) survivor is retained. The number of survivors never exceeds the number of states in the trellis.

The basic operations required in a soft decision Viterbi decoder are:[Ref 3:p 651]

- computation of branch metrics, for an additive white gaussian noise channel, are proportional to $(r-x_i)^2$ where r is the received sample and x_i is the noise free symbol associated with the message.

- addition of branch metrics and survivors to determine the survivors for each state.

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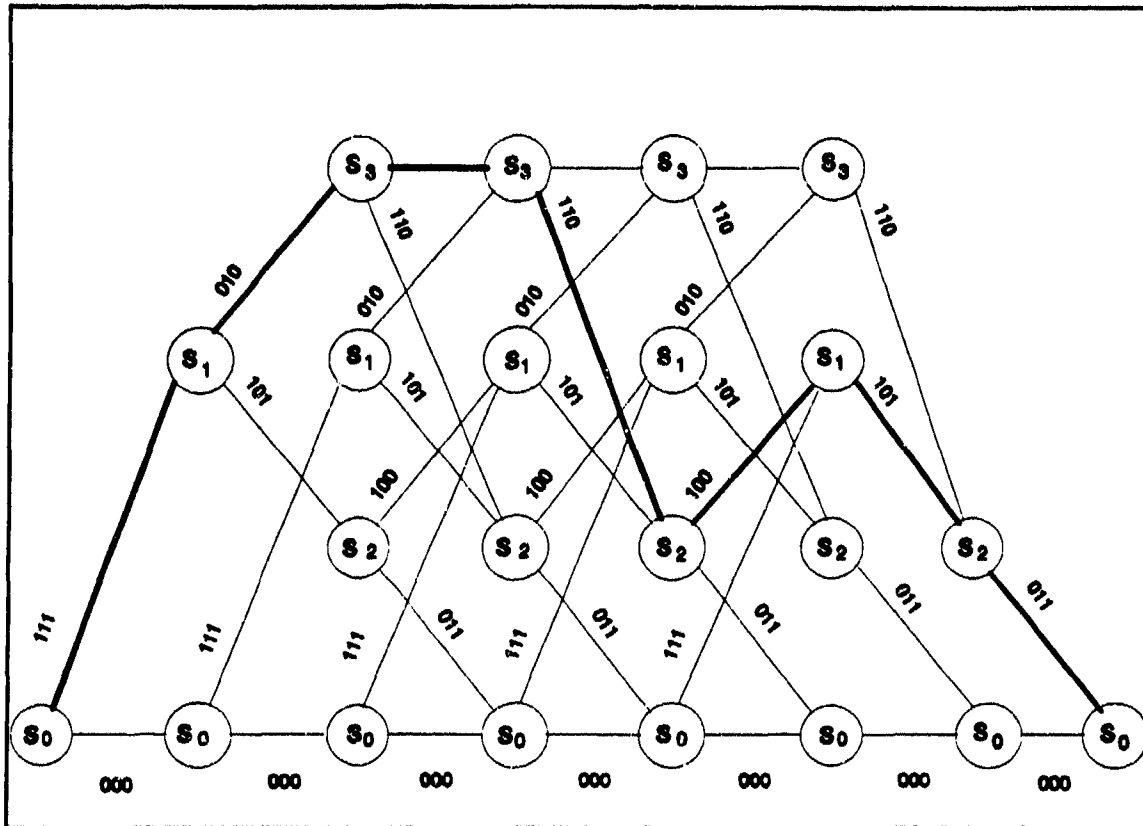
Figure 4 Trellis diagram for a (3,1,2) code. (from Ref 4:p 316)

- comparison along the extended path metrics to determine the survivor for each state.

It was shown that the Viterbi algorithm was equivalent to a dynamic programming solution to the problem of finding the shortest path through a weighted graph [Ref 3:pp 317-321]. The decoder must produce an estimate $\hat{v}$ of the codeword v based on the received sequence r. For an information sequence of length L, the trellis must contain L+m+1 levels or time units to decode the sequence. An (n,k,m) code has an information sequence of length K1, and is encoded into a codeword of length N=n(L+m). A maximum likelihood decoder (MLD) for a discrete memoryless channel (DMC) chooses

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$\hat{v}$ as the codeword v which maximizes the log-likelihood function log $P(r|v)$. Since for a DMC:

$$P(r \mid v) = \prod_{i=0}^{L+m-1} P(r_i \mid v_i) = \prod_{i=0}^{N-1} P(r_i \mid v_i) \qquad (7)$$

it follows that log-likelihood function is formed by the summation of the branch metrics:

$$\log P(r \mid v) = \sum_{i=0}^{L+m-1} \log P(r_i \mid v_i) = \sum_{i=0}^{N-1} \log P(r_i \mid v_i) \qquad (8)$$

where $P(r_i|v_i)$ is a channel transition probability. This is a minimum error probability decoding rule when all code words are equally likely. The log-likelihood function log $P(r|v)$ are called the metric of path v and is denoted $M(r|v)$. A partial path metric is formed by summing up the partial path metrics for j branches and is expressed by:

$$M([r \mid v]_j) = \sum_{i=0}^{j-1} M(r_i \mid v_i) \qquad (9)$$

The final survivor $\hat{v}$ in the Viterbi algorithm is the maximum likelihood path;

$$M(r \mid \hat{v}) \geq M(r \mid v), \qquad \text{all } v \neq \hat{v} \qquad (10)$$

The path is then traced back to determine the symbol transmitted.[Ref. 4:pp 316-318]

III. SYSTEM DEVELOPMENT

All the programs used in this thesis were run on an Ariel PC-56 DSP board. The PC-56 DSP board is an eight bit card that can be used in any industry standard computer. The primary components of the PC-56 are a 20 Mhz Motorola DSP 56001, 16k of external memory, a 24 bit bidirectional interrupt-driven port with a header for external I/O, a TLC32040 14 bit analog interface chip with built in input and output filters gain section. An Ariel PC-56D was also used for comparison, it contains a 27 Mhz Motorola DSP 56001, 64k of external memory and a NEXT compatible DB-15 port. For actual implementation of this system a DSP board with dual A/D and D/A converters is required.

The structure of the Motorola DSP 56001 makes it ideally suited for high speed communications. In the DSP chip there are 256 locations of 24 bit x and y memory that occupy the lowest 256 locations of the DSP address space. The locations from 256-511 are allocated for the on-chip ROM. The onboard ROM contains Mu-law and A-law expansion tables as well as a full four quadrant sine table. There are 512 locations of 24 bit high speed program RAM (PRAM) on the chip.[Ref 5]

A feature that makes the DSP 56001 desirable for use in this system is that it offers several addressing modes. It implements three types of arithmetic for addressing, linear, modulo and reverse carry. For each of the address registers R0-R7 there is an offset register N0-N7 and a modifier register M0-M7. The offset register contains the values to increment and decrement the address register. The modifier register

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defines the type of address arithmetic to be used. For modulo arithmetic the contents of the modifier register Mn specify the base modulus.[Ref 6:pp 5-2,5-4] The DSP 56001 also provides three different addressing modes, register direct, address register indirect, and post or pre increment/decrement.

The basic configuration of an MFM transmitter and receiver are shown in Figure 5 and Figure 6. For the purpose of this project three types of modulators were used in the investigation for use on a DSP board. In previous systems implemented at NPS [Ref 1] and [Ref 7] the complex conjugate of the input data was loaded into the image frequencies of the IFFT. This resulted in only real data being generated by the IFFT. For this project complex data was loaded into all available locations to generate both real and complex data. The signal may be modulated using any type of modulation scheme. For the purpose of this thesis the three different modulators used a 16 QAM signal, a 2/3 rate convolutional encoder with an eight PSK signal and a 4/5 rate code with a 32 QAM signal using the CCITT modem standard V.32 convolutional encoder. The source code for the V.32 encoder and Viterbi decoder was included in an example manual from Motorola [Ref 8:pp A1-C2] for use on their simulator and was modified for use on the Ariel board to generate a Multi-Frequency Modulated signal.

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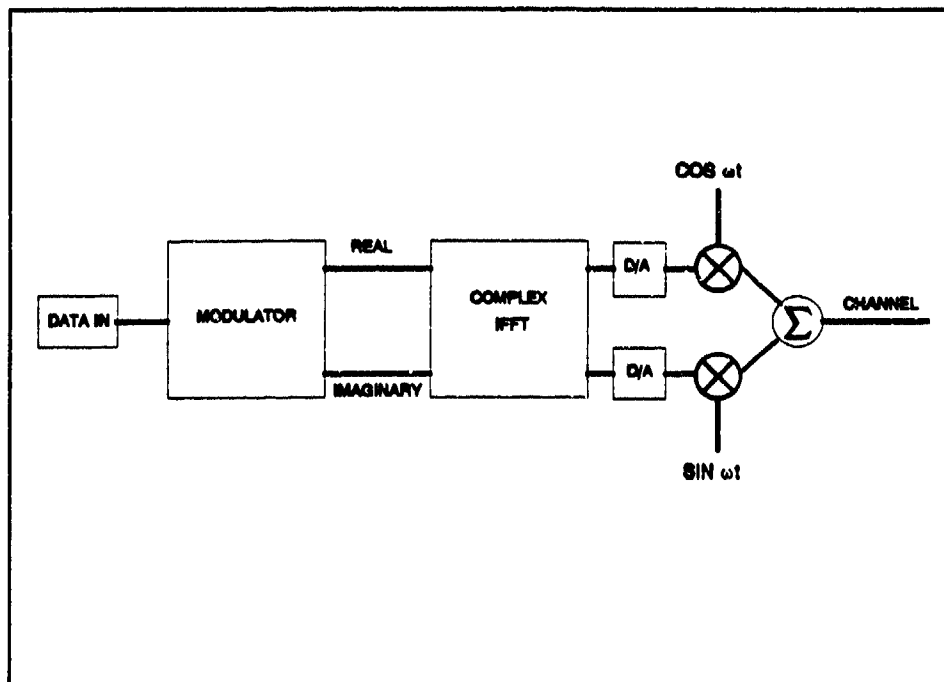
Figure 5 MFM Transmitter
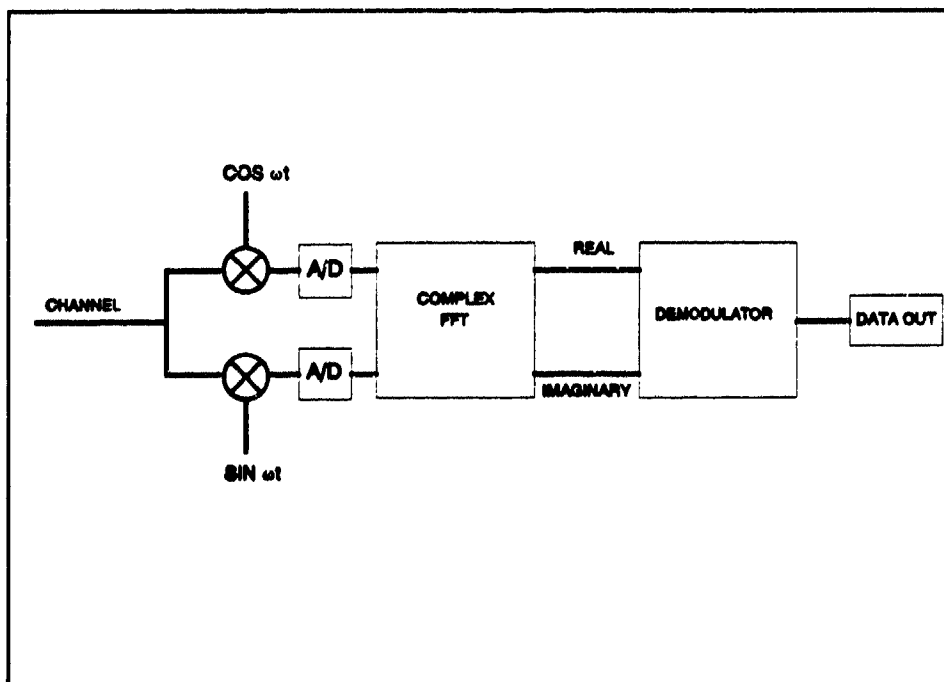
Figure 6 MFM Receiver

A. ENCODERS

Of the 256 bins available in the IFFT array, only 201 message symbols are carried in each baud. The other 55 bins are loaded with zeros to allow for filtering of the signal. Once the 256 bins are filled, the IFFT subroutine is called to generate the values for the MFM signal. At this point the values would be clocked out through the D/A converters.

1. 16 QAM Encoder

At initialization of the system the message to be transmitted is downloaded from the host computer as well as the programs for encoding (QAM16EN) the signal and the data for the look-up table (16QAMRE.DAT and 16QAMIM.DAT). The data is read in using four bit increments. The value read in is moved into an offset register and used to determine the coordinates of the points on the constellation in Figure 7. The real values of the constellation are stored in the x memory and the corresponding imaginary values are stored in the y memory. Once the constellation values are determined the real and complex values are stored in memory. Once one hundred samples have been read into memory the routine installs 55 zeros at the center of the array then reads the remainder of the 256 symbols. Once all 256 samples are stored the IFFT routine is called to generate the time domain signal.

2. 2/3 Rate Convolutional Encoder

At initialization of the system the message to be sent as well as the encoder (23ENCOD) and the data files (23REAL.DAT and 23IMAG.DAT) are loaded into memory. In the 2/3 rate encoder two bits are read in from the input message and using the convolutional encoder of Figure 2 a third bit
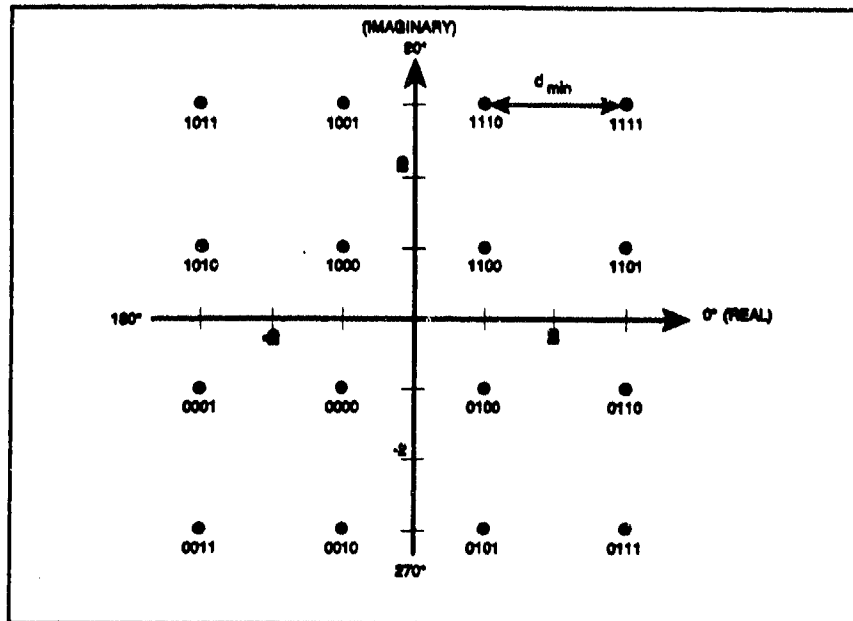
Figure 7 CONSTELLATION FOR 16QAM SIGNAL (from Ref.8:p 2-2)

is generated. The eight PSK constellation in Figure 8 is used to transmit the message. A look up table is used to find the points on the constellation. After the points on the constellation are found, the real
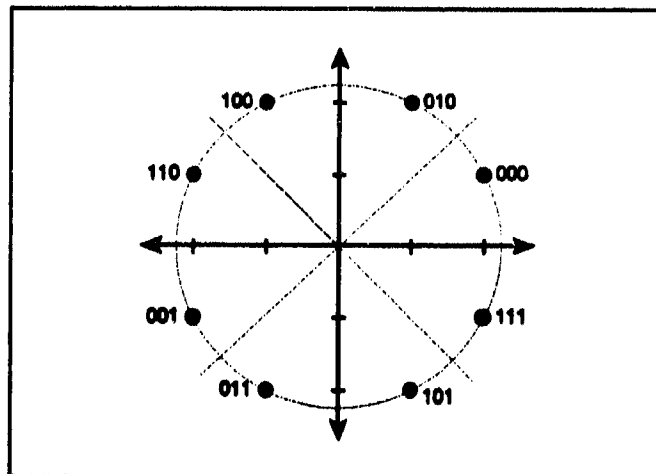
Figure 8   2/3   Rate   Code Constellation

and imaginary values are loaded into an array. After the first 100 samples are loaded into the array, 55 zeros are loaded into the array to allow for filtering and the remainder of symbols are read. Once all 256 values are loaded the IFFT is performed and the values are stored ready to be clocked out.

3. V.32 Convolutional Encoder

At initialization the message, encoder program (MFMENCOD) and data files (QAMREAL.DAT and QAMIMAG.DAT) are downloaded to the to the DSP board. For the V.32 encoder after the four bit symbol is read, the convolutional encoder Figure 9 is used to generate a fifth redundant bit.
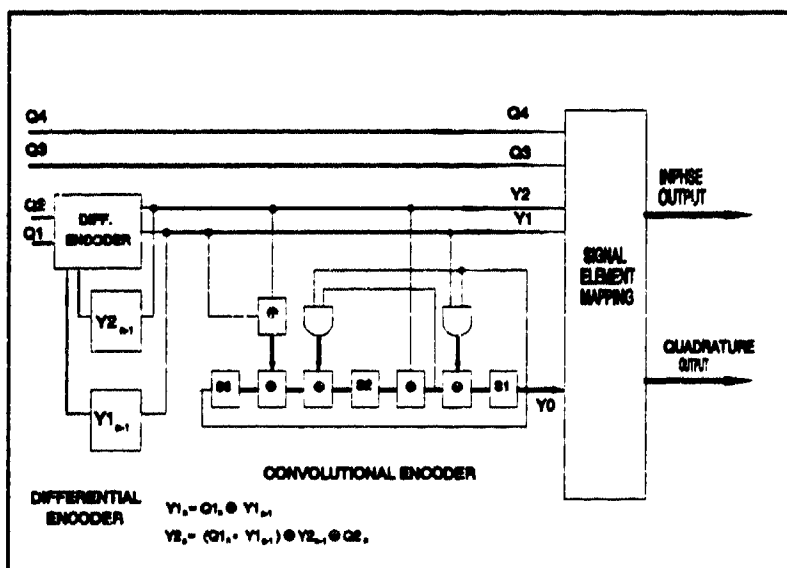
Figure 9 V.32 Convolutional Encoder (from Ref. 8:p 2-4)

The 32 QAM constellation in Figure 10 is used for the five bit signal. A look up table is used to determine the real and imaginary values for the corresponding point on the constellation. As in the 16 QAM encoder and

the 2/3 rate code encoder, 201 encoded symbols and 55 zeros are loaded into the IFFT array and the IFFT is performed.

B. GENERATING THE MFM SIGNAL

When all 256 bins have been loaded the program calls the routine to perform the IFFT. The routine used to perform the IFFT was generated by using the relationship $x(n)=(FFT[X^t(k)]^t/N$. The program to generate the FFT was included with the Ariel DSP package[1]. In order to speed up
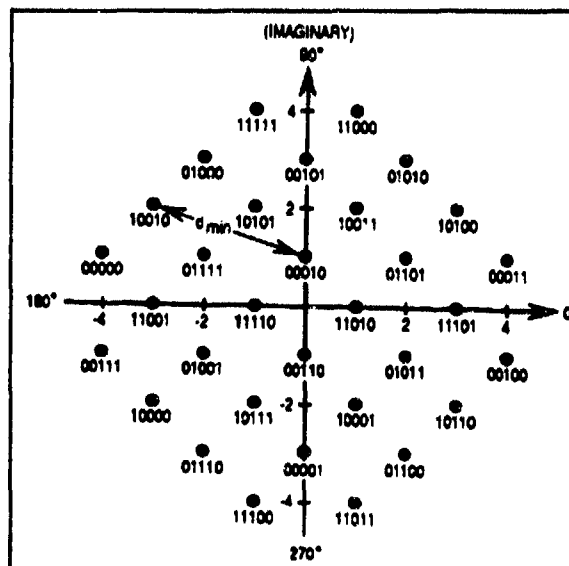
Figure 10 32 QAM Constellation
(from Ref. 8:p 2-3)

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the routine as much as possible the values used in the look up table were conjugated and scaled by 256 to save from doing them in the modulator. By using the prescaled and conjugated values in the encoder constellations the IFFT routines require no more time than the FFT routine used in the decoder.

After computing the IFFT the real and complex values are stored in memory. By utilizing a DSP board with dual D/A converters the real and imaginary values can be clocked out at the sampling frequency f_χ. The design of the system may be altered to either allow transmission of each 256 symbol packet or store the entire message and transmit at the end of the message.

C. DECODERS

After using a quadrature receiver to recover the inphase and quadrature components, the signals are sampled at the same frequency, f_χ, as used in the encoder. These sampled values are used to determine the transmitted symbols.

1. 16 QAM DECODER

At initialization the boundary data file BOUN16.D and the decoder program QAM16DEC must be loaded onto the DSP board. To decode the 16 QAM signal the constellation was partitioned as shown in Figure 11. The magnitude of each received signal value is compared to the boundaries in the first quadrant. Once a bounded area is found the actual values are then used to determine the correct quadrant. The data file BOUN16.D is indexed and the corresponding symbol is read out.
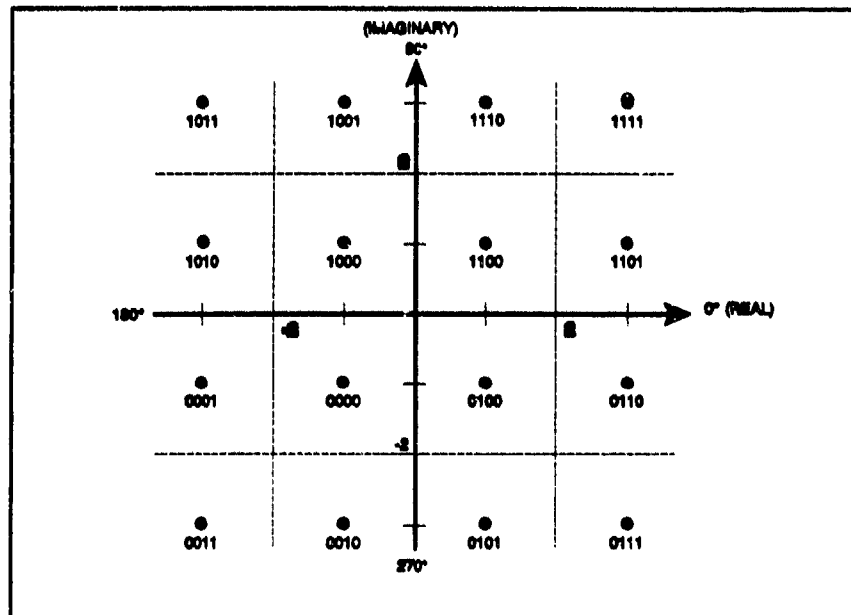
Figure 11 PARTITIONED 16 QAM CONSTELLATION

2.   2/3 Rate Code Decoder

The use of a convolutional encoder requires the use of a soft decoder such as a Viterbi Decoder to recover the transmitted symbol. The boundary data file BOUN23.D and the Viterbi decoder program MFM23DEC are first loaded into memory and run. After computing the FFT to recover the transmitted symbols, a received point is read out of the array and the quadrant of the constellation (Figure 8) that it lies in is found. The data file BOUN23.D holds the four closest points, one in each state, to the quadrant that the received point lies in. A boundary file is used to help speed up the decoder routine, instead of computing the distance to all points in the constellation only the point closest in each state is used for computing the euclidean distance to the received point. Using the

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received point and the four values found in BOUN23.D the euclidean distance to each point is computed. The accumulated distance table is updated with the distance to each state. The smallest accumulated distance is found and the trellis is traced back 16 time periods to find the most likely transmitted point. After finding the most likely point it must be decoded. The bit Y2 is discarded and using the relation Y1=X1⊕S1 the value of X1 is recovered by using past state information and X1=Y1⊕S1. The value of Y0 is output as X0. The values are stored in memory.

3. V.32 DECODER

The use of a convolutional encoder in the V.32 system requires that a Viterbi Decoder must also be used to recover the transmitted symbol. The boundary file BOUND.D and the program file MFMDECOD are first loaded into memory. The constellation must first be partitioned in such a way that the partition equally divides the distance between the four symbols in the same state. The partition used for the state 110 is shown in Figure 12. By using this method the partitions for all eight states may be imposed on the constellation resulting in 52 separate partitions as shown in Figure 13. For each partition there are eight points, one from each state, that is closest to the boundary. These eight points for each of the 52 boundaries are stored in the data file BOUND.D that is loaded into the processor at initialization. Once a bounded area has been found, the eight points are used to determine the euclidean distance from the received point to the points read from BOUND.D. The distances calculated
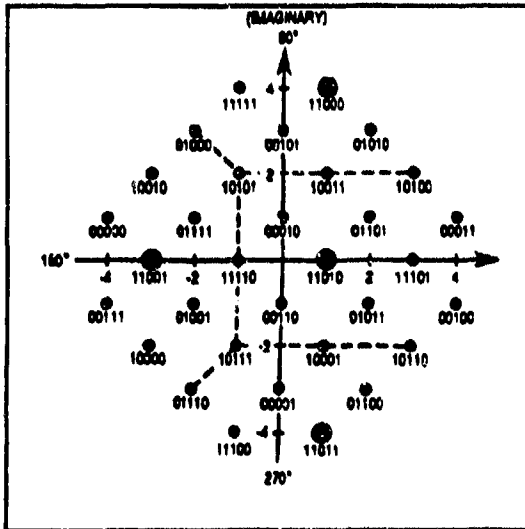
Figure 12 32 QAM Constellation (from Ref. 8: p 3-7)
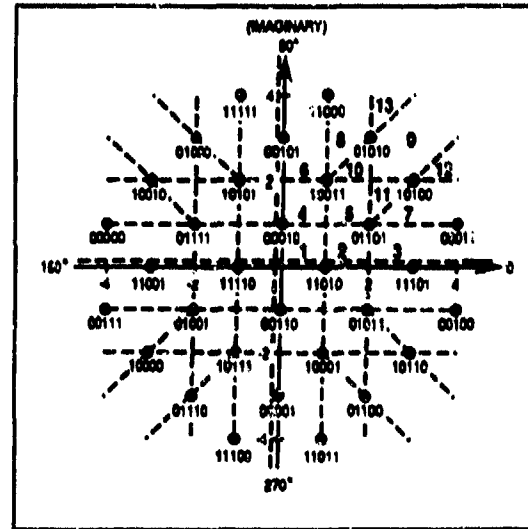
Figure 13 Partitioned 32 QAM Constellation (from Ref. 8:p 3-7)

are used to update the accumulated distance to each state. The state with the smallest accumulated distance is found and the trellis is traced back 16 time periods to find the input state at the end of the most likely path. The closest point in that state to the input point at that time period is found and output. The most significant bit is masked off. The differential encoding used on the two most significant bits must be decoded by:

$$Q1_n = Y1_n \oplus Y1_{n-1} \qquad (11)$$

$$Q2_n = (Q1_n \cdot Y1_{n-1}) \oplus Y2_{2n-1} \oplus Y2_n \qquad (12)$$

Once the two most significant bits have been decoded the symbol is output to memory.

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IV. SYSTEM OPERATION

The Digital Signal Processing boards used to run the programs that were developed did not have dual D/A, A/D converters required to fully implement an MFM system. The programs as listed, do not fully implement a MFM system. There are no subroutines included for the clocking out of samples through D/A converters in the encoders and no sampling routines in the decoders. The encoding programs read in the data to be transmitted from memory, encode the data and store the values created by the IFFT in memory. The same locations are used to store all values generated, it is assumed that the values will be clocked out once available freeing up these locations for the next baud. The values would be clocked out through the D/A converters to generate a signal. In the decoder programs, the data is read out of locations were the sampled values would be stored. The same locations are used for all sampled values, it is assumed each baud will be decoded prior to receiving the next baud. If a convolutional encoder is used, a decision must be made as to whether the next baud should be delayed to allow decoding of the previous baud, or transmit the baud when available and store the values in the decoder for off line decoding.

Once the modulation technique is chosen the appropriate constellation data files must be loaded for the look up tables used in the decoder. These files are located in Appendices D, H, L. Once the appropriate data files are loaded the encoder program and message file are loaded into memory and the encoder is run. The programs developed were run out of

Ariel's BUG-56 monitor-debugger, they can also be adapted to run as subroutines in a Microsoft C program. The monitor-debugger allowed easy access to all the chips registers and functions making running and debugging the programs simpler.

To use the decoder the appropriate boundary values must be loaded. Once the boundary values are loaded the decoder program is loaded and run. The output is stored in memory.

All of the programs listed are written in assembly language for the DSP 56001 and must be compiled into a loadable file using the Motorola assembler. All of the data files are properly formatted to be read into the correct memory locations. If the programs are run out of the BUG-56 monitor-debugger, a macro can be created to load all the files needed for an encoder or decoder using one instruction.

V. CONCLUSIONS

The time required for the encoder/decoder for a 256 sample baud is

shown in Table I. The Turbo Pascal routine used in past projects

Table I Sample Times for Programs Using 256 Sample Baud

PROGRAMS	TIME(ms)	
	20 Mhz 56001	27 Mhz 56001
4096 pt complex fft	.023 sec	.017 sec
256 pt complex fft	1	.7
MFM PROGRAMS (for 256 sample baud)		
16 QAM Encoder	2.7	2.0
16 QAM Decoder	2.7	2.0
2/3 convolutional encoder	3.2	2.4
2/3 rate code Viterbi decoder	10.8	8.0
V.32 encoder (w/o diff enc)	3.6	2.7
V.32 Viterbi decoder (w/o diff enc)	15.8	11.7
V.32 encoder	3.7	2.8
V.32 Viterbi decoder	16.9	12.5

required 22 seconds to perform a 4096 point complex FFT on an 8 Mhz AT computer with a coprocessor.[Ref. 9:p 18] It is readily apparent that the Digital Signal Processor has a tremendous performance advantage over normal software routine. Using Matlab benchmarks a Sun Sparc station required .34 seconds and a 20 Mhz 386 with a coprocessor required 1.16 seconds for a 4096 point complex FFT.

Multi-Frequency Modulation can be easily implemented on a Digital Signal Processor Board. The programs for implementing the 16 QAM system required the same amount of time to encode and to decode. There is little penalty for using a trellis encoder but the Viterbi decoder takes over three times longer than the encoder for the 2/3 rate code and almost four times as long for the V.32 decoder.

The use of DSP boards in implementing MFM shows much improvement over existing techniques. Suggestions for future research include running the MFM programs on DSP boards with the required dual D/A, A/D converters. A determination must be made as to whether the decoder should store the entire message or delay transmission of a baud until the previous baud is decoded. Another alternative is to use more than one processor, one to encode/decode and another to clock out/sample the data.

The versatility and speed of Digital Signal Processors make them ideally suited for use for a Multi-Frequency Modulated System. They offer much greater performance than is currently available using software routines and are can easily perform other functions by simply loading a new program. The continuing development of more DSP compatible products show that there will be greater uses of DSP technology in the future.

```
;
; This program originally available on the Motorola DSP bulletin board.
; It is provided under a DISCLAIMER OF WARRANTY available from
; Motorola DSP Operation, 6501 Wm. Cannon Drive W., Austin, Tx., 78735.
;
; Radix 2, In-Place, Decimation-In-Time FFT (fast).
;
; Last Update 18 Aug 88    Version 1.0
;
fft   macro   points,data,odata,coef
fft   ident   1,0
;
; Radix 2 Decimation in Time In-Place Fast Fourier Transform Routine
;
;         Complex input and output data
;         Real data in X memory
;         Imaginary data in Y memory
;         Normally ordered input data
;         Normally ordered output data
;         Coefficient lookup table
;         -Cosine values in X memory
;         -Sine values in Y memory
;
; Macro Call - fft    points,data,odata,coef
;
;         points      number of points (16-32768, power of 2)
;         data        start of data buffer
;         odata       start of output data buffer
;         coef        start of sine/cosine table
;
; Alters Data ALU Registers
;         x1      x0      y1      y0
;         a2      a1      a0      a
;         b2      b1      b0      b
;
; Alters Address Registers
;         r0      n0      m0
;         r1      n1      m1
;                 n2
;
;         r4      n4      m4
;         r5      n5      m5
;         r6      n6      m6
;
```

```
; Alters Program Control Registers
;         pc      sr
;
; Uses 6 locations on System Stack
;
; Latest Revision - 18 Aug-88
;

      move #data,r0              ;initialize input pointer
      move #points/4,n0          ;initialize input and output pointers offset
      move n0,n4                 ;
      move n0,n6                 ;initialize coefficient offset
      move #points-1,m0          ;initialize address modifiers
      move m0,m1                 ;for modulo addressing
      move m0,m4
      move m0,m5
;
; Do first and second Radix 2 FFT passes, combined as 4-point butterflies
;
      move           x:(r0)+n0,x0
      tfr  x0,a      x:(r0)+n0,y1

      do   n0,_twopass
      tfr  y1,b      x:(r0)+n0,y0
      add  y0,a      x:(r0),x1                      ;ar+cr
      add  x1,b      r0,r4                          ;br+dr
      add  a,b       (r0)+n0                        ;ar'=(ar+cr)+(br+dr)
      subl b,a       b,x:(r0)+n0                    ;br'=(ar+cr)-(br+dr)
      tfr  x0,a      ,x0           y:(r0),b
      sub  y0,a                    y:(r4)+n4,y0     ;ar-cr
      sub  y0,b      x0,x:(r0)                      ;bi-di
      add  a,b                     y:(r0)+n0,x0     ;cr'=(ar-cr)+(bi-di)
      subl b,a       b,x:(r0)                       ;dr'=(ar-cr)-(bi-di)
      tfr  x0,a      ,x0           y:(r4),b
      add  y0,a                    y:(r0)+n0,y0     ;bi+di
      add  y0,b      x0,x:(r0)+n0                   ;ai+ci
      add  b,a                     y:(r0)+,x0       ;ai'=(ai+ci)+(bi+di)
      subl a,b                     a,y:(r4)+n4      ;bi'=(ai+ci)-(bi+di)
      tfr  x0,a                    b,y:(r4)+n4
      sub  y0,a      x1,b                           ;ai-ci
      sub  y1,b      x:(r0)+n0,x0                   ;dr-br
      add  a,b       x:(r0)+n0,y1                   ;ci'=(ai-ci)+(dr-br)
      subl b,a                     b,y:(r4)+n4      ;di'=(ai-ci)-(dr-br)
      tfr  x0,a                    ,y:(r4)+
_twopass
;
; Perform all next FFT passes except last pass with triple nested DO loop
;
      move #points/8,n1          ;initialize butterflies per group
      move #4,n2                 ;initialize groups per pass
      move #-1,m2                ;linear addressing for r2
```

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```
        move  #0,m6                    ;initialize C address modifier for
                                       ;reverse carry (bit-reversed) addressing


        do    #@cvi(@log(points)/@log(2)-2.5),_end_pass    ;example: 7 passes
for 1024 pt. FFT
        move  #data,r0                              ;initialize A input
pointer
        move  r0,r1
        move  n1,r2
        move  r0,r4                                 ;initialize A
output pointer
        move  (r1)+n1                               ;initialize B input
pointer
        move  r1,r5                                 ;initialize B
output pointer
        move  #coef,r6                              ;initialize C input
pointer
        lua   (r2)+,n0                              ;initialize pointer
offsets
        move  n0,n4
        move  n0,n5
        move  (r2)-                                 ;butterfly loop
count
        move          x:(r1),x1      y:(r6),y0      ;lookup -sine and
-cosine values
        move          x:(r6)+n6,x0   y:(r0),b       ;update C pointer,
preload data
        mac   x1,y0,b                y:(r1)+,y1
        macr  -x0,y1,b               y:(r0),a


        do    n2,_end_grp
        do    r2,_end_bfy
        subl  b,a        x:(r0),b        ,y:(r4)        ;Radix 2 DIT
butterfly kernel
        mac   -x1,x0,b   x:(r0)+,a       ,y:(r5)
        macr  -y1,y0,b   x:(r1),x1
        subl  b,a        b,x:(r4)+       y:(r0),b
        mac   x1,y0,b                    y:(r1)+,y1
        macr  -x0,y1,b   a,x:(r5)+       y:(r0),a
_end_bfy
        move  (r1)+n1
        subl  b,a        x:(r0),b        ,y:(r4)
        mac   -x1,x0,b   x:(r0)+n0,a     ,y:(r5)
        macr  -y1,y0,b   x:(r1),x1       y:(r6),y0
        subl  b,a        b,x:(r4)+n4     y:(r0),b
        mac   x1,y0,b    x:(r6)+n6,x0    y:(r1)+,y1
        macr  -x0,y1,b   a,x:(r5)+n5     y:(r0),a
_end_grp
        move  n1,b1
        lsr   b     n2,a1      ;divide butterflies per group by two
        lsl   a     b1,n1      ;multiply groups per pass by two
```

```
        move al,n2
_end_pass
;
; Do last FFT pass
;
        move #2,n0              ;initialize pointer offsets
        move n0,n1
        move #points/4,n4       ;output pointer A offset
        move n4,n5              ;output pointer B offset
        move #data,r0           ;initialize A input pointer
        move #odata,r4          ;initialize A output pointer
        move r4,r2              ;save A output pointer
        lua  (r0)+,r1           ;initialize B input pointer
        lua  (r2)+n2,r5         ;initialize B output pointer
        move #0,m4              ;bit-reversed addressing for output ptr. A
        move m4,m5              ;bit-reversed addressing for output ptr. B
        move #coef,r6           ;initialize C input pointer
        move (r5)-n5            ;predecrement output pointer
        move            x:(r1),x1      y:(r6),y0
        move            x:(r5),a       y:(r0),b

        do   n2,_lastpass
        mac  x1,y0,b   x:(r6)+n6,x0    y:(r1)+n1,y1   ;Radix 2 DIT butterfly
kernel
        macr -x0,y1,b  a,x:(r5)+n5     y:(r0),a       ;with one butterfly per
group
        subl b,a       x:(r0),b        b,y:(r4)
        mac  -x1,x0,b  x:(r0)+n0,a     a,y:(r5)
        macr -y1,y0,b  x:(r1),x1       y:(r6),y0
        subl b,a       b,x:(r4)+n4     y:(r0),b
_lastpass
        move            a,x:(r5)+n5
        endm
```

; This program uses the same routine used in the program FFT. The IFFT is
; calculated using the relationship IFFT=(FFT[X(k)]/N). It is assumed that
; the values used in the constellation are conjugated and prescaled by N.
; If the values to be used are not conjugated a routine must be added to
; the beginning of this routine to conjugate the values. The division by
; N can also be accomplised by using left shifts, but will slow down the
; program.
;
; Radix 2, In-Place, Decimation-In-Time IFFT (fast).
;
;
ifft   macro   points,data,odata,coef
ifft   ident   1,0
;
; Radix 2 Decimation in Time In-Place Inverse Fast Fourier Transform
; Routine
;
;   Complex input and output data
;   Real data in X memory
;   Conjugated and prescaled imaginary data in Y memory
;   Normally ordered input data
;   Normally ordered output data
;   Coefficient lookup table
;   -Cosine values in X memory
;   -Sine values in Y memory
;
; Macro Call - ifft   points,data,odata,coef
;
;   points   number of points (16-32768, power of 2)
;   data   start of data buffer
;   odata   start of output data buffer
;   coef   start of sine/cosine table
;
; Alters Data ALU Registers
;   x1   x0   y1   y0
;   a2   a1   a0   a
;   b2   b1   b0   b
;
; Alters Address Registers
;   r0   n0   m0
;   r1   n1   m1
;   n2
;
;   r4   n4   m4
;   r5   n5   m5

```
;         r6      n6      m6
;
; Alters Program Control Registers
;       pc        sr
;
; Uses 6 locations on System Stack
;
; Latest Revision - 18 Aug-88
;
    move  #data,r0              ;initialize input pointer
    move  #points/4,n0          ;initialize input and output pointers offset
    move  n0,n4                 ;
    move  n0,n6                 ;initialize coefficient offset
    move  #points-1,m0          ;initialize address modifiers
    move  m0,m1                 ;for modulo addressing
    move  m0,m4
    move  m0,m5


;
; Do first and second Radix 2 IFFT passes, combined as 4-point butterflies
;

    move            x:(r0)+n0,x0
    tfr   x0,a      x:(r0)+n0,y1

    do    n0,_twopass
    tfr   y1,b      x:(r0)+n0,y0
    add   y0,a      x:(r0),x1                       ;ar+cr
    add   x1,b      r0,r4                           ;br+dr
    add   a,b       (r0)+n0                         ;ar'=(ar+cr)+(br+dr)
    subl  b,a       b,x:(r0)+n0                     ;br'=(ar+cr)-(br+dr)
    tfr   x0,a      a,x0          y:(r0),b
    sub   y0,a                    y:(r4)+n4,y0      ;ar-cr
    sub   y0,b      x0,x:(r0)                       ;bi-di
    add   a,b                     y:(r0)+n0,x0      ;cr'=(ar-cr)+(bi-di)
    subl  b,a       b,x:(r0)                        ;dr'=(ar-cr)-(bi-di)
    tfr   x0,a      a,x0          y:(r4),b
    add   y0,a                    y:(r0)+n0,y0      ;bi+di
    add   y0,b      x0,x:(r0)+n0                    ;ai+ci
    add   b,a                     y:(r0)+,x0        ;ai'=(ai+ci)+(bi+di)
    subl  a,b                     a,y:(r4)+n4       ;bi'=(ai+ci)-(bi+di)
    tfr   x0,a                    b,y:(r4)+n4
    sub   y0,a      x1,b                            ;ai-ci
    sub   y1,b      x:(r0)+n0,x0                    ;dr-br
    add   a,b       x:(r0)+n0,y1                    ;ci'=(ai-ci)+(dr-br)
    subl  b,a                     b,y:(r4)+n4       ;di'=(ai-ci)-(dr-br)
    tfr   x0,a                    a,y:(r4)+
_twopass
;
; Perform all next IFFT passes except last pass with triple nested DO loop
;
```

```
        move #points/8,n1      ;initialize butterflies per group
        move #4,n2             ;initialize groups per pass
        move #-1,m2            ;linear addressing for r2
        move #0,m6            ;initialize C address modifier for
                             ;reverse carry (bit-reversed) addressing

        do   #@cvi(@log(points)/@log(2)-2.5),_end_pass   ;example: 7 passes
for 1024 pt. IFFT
        move #data,r0                                ;initialize A input
pointer
        move r0,r1
        move n1,r2
        move r0,r4                                   ;initialize A output
pointer
        move (r1)+n1                                 ;initialize B input
pointer
        move r1,r5                                   ;initialize B output
pointer
        move #coef,r6                                ;initialize C input
pointer
        lua  (r2)+,n0                                ;initialize pointer
offsets
        move n0,n4
        move n0,n5
        move (r2)-                                   ;butterfly loop count
        move          x:(r1),x1      y:(r6),y0       ;lookup -sine and -cosine
values
        move          x:(r6)+n6,x0   y:(r0),b        ;update C pointer, preload
data
        mac x1,y0,b                  y:(r1)+,y1
        macr -x0,y1,b                y:(r0),a


        do   n2,_end_grp
        do   r2,_end_bfy
        subl b,a       x:(r0),b      b,y:(r4)        ;Radix 2 DIT butterfly
kernel
        mac  -x1,x0,b  x:(r0)+,a     a,y:(r5)
        macr -y1,y0,b  x:(r1),x1
        subl b,a       b,x:(r4)+     y:(r0),b
        mac  x1,y0,b                 y:(r1)+,y1
        macr -x0,y1,b  a,x:(r5)+     y:(r0),a
_end_bfy
        move (r1)+n1
        subl b,a       x:(r0`,b      b,y:(r4)
        mac  -x1,x0,b  x:(r0)+n0,a   a,y:(r5)
        macr -y1,y0,b  x:(r1),x1     y:(r6),y0
        subl b,a       b,x:(r4)+n4   y:(r0),b
        mac  x1,y0,b   x:(r6)+n6,x0  y:(r1)+,y1
        macr -x0,y1,b  a,x:(r5)+n5   y:(r0),a
_end_grp
        move n1,b1
```

```
        lsr   b    n2,a1       ;divide butterflies per group by two
        lsl   a    b1,n1       ;multiply groups per pass by two
        move a1,n2
_end_pass
;
; Do last IFFT pass
;
        move #2,n0             ;initialize pointer offsets
        move n0,n1
        move #points/4,n4      ;output pointer A offset
        move n4,n5             ;output pointer B offset
        move #data,r0          ;initialize A input pointer
        move #odata,r4         ;initialize A output pointer
        move r4,r2             ;save A output pointer
        lua  (r0)+,r1          ;initialize B input pointer
        lua  (r2)+n2,r5        ;initialize B output pointer
        move #0,m4             ;bit-reversed addressing for output ptr. A
        move m4,m5             ;bit-reversed addressing for output ptr. B
        move #coef,r6          ;initialize C input pointer
        move (r5)-n5           ;predecrement output pointer
        move            x:(r1),x1      y:(r6),y0
        move            x:(r5),a       y:(r0),b


        do   n2,_lastpass
        mac  x1,y0,b  x:(r6)+n6,x0   y:(r1)+n1,y1   ;Radix 2 DIT butterfly
kernel
        macr -x0,y1,b  a,x:(r5)+n5   y:(r0),a       ;with one butterfly per
group

        subl b,a                                    ;complete last butterfly
        neg  b                                       ;and output conjugate
        move b,y:(r4)
        move x:(r0),b
        neg  a
        move a,y:(r5)
        mac  -x1,x0,b  x:(r0)+n0,a
        macr -y1,y0,b  x:(r1),x1       y:(r6),y0
        subl b,a       b,x:(r4)+n4     y:(r0),b
_lastpass
        move            a,x:(r5)+n5
        endm
```

APPENDIX C.   SINE COSINE GENERATOR

```
;
; This program originally available on the Motorola DSP bulletin board.
; It is provided under a DISCLAIMER OF WARRANTY available from
; Motorola DSP Operation, 6501 Wm. Cannon Drive W., Austin, Tx., 78735.
;
; Sine-Cosine Table Generator for FFTs.
;
; Last Update 25 Nov 86   Version 1.2
;
sincos   macro   points,coef
sincos   ident   1,2
;
;        sincos  -       macro to generate sine and cosine coefficient
;                        lookup tables for Decimation in Time FFT
;                        twiddle factors.
;
;        points  -       number of points (2 - 32768, power of 2)
;        coef    -       base address of sine/cosine table
;                        negative cosine value in X memory
;                        negative sine value in Y memory
;
; Latest revision - 25-Nov-86
;

pi       equ     3.141592654
freq     equ     2.0*pi/@cvf(points)

         org     x:coef
count    set     0
         dup     points/2
         dc      -@cos(@cvf(count)*freq)
count    set     count+1
         endm

         org     y:coef
count    set     0
         dup     points/2
         dc      -@sin(@cvf(count)*freq)
count    set     count+1
         endm

         endm    ;end of sincos macro
```

APPENDIX D.   16 QAM DATA FILES

The following two data files are prescaled (full scale constellation values divided by 256) and conjugated. Both data files must be loaded into the DSP prior to running QAM16ENC. They are properly formatted to be read into memory. They occupy the first 16 locations of x and y memory.

Data file '16QAMRE.DAT'

```
HX
00000000,0000000F
FFF800,FFE800,FFF800,FFE800
000800,000800,001800,001800
FFF800,FFF800,FFE800,FFE800
000800,001800,000800,001800
```


Data file '16QAMIM.DAT'

```
HY
00000000,0000000F
000800,000800,001800,001800
000800,001800,000800,001800
FFF800,FFE800,FFF800,FFE800
FFF800,FFF800,FFE800,FFE800
```

```
; This program is an encoder for a 16 QAM system.  A symbol consisting of
; four bits is read into the engwder and using a look up table outputs the
; values for the points on the constellation and are loaded into an array.
; The data files 16QAMRE.DAT and 16QAMIM.DAT must be loaded prior to
; running 16QAMEN.  Once the 256 values are loaded into an array an IFFT
; is performed to generate the MFM signal.
;
; The following registers are modified:
;    r2
;    r3
;    r4
;    r5
;
; The following Data ALU Registers are modified:
;    a0   a1   a
;    x0   x1   y0
;    b0   b1   b
;


qam16en


        ident   1,1
    page    132,54
    opt     nomd,nomex,loc

    include   'sincos'  ;include macro for sine cosine values
    include   'ifft'    ;include macro for calculating IFFT

; Define memory locations to be used in the program

start    equ        $100   ;starting program location
startifft    equ  $350   ;starting program location for IFFT
input    equ     40      ;starting location of input bits
output   equ     50      ;starting location of output bits
coef     equ     128     ;location of coefficients for IFFT
points   equ     256     ;number of points for IFFT
data     equ     768     ;location of input data for IFFT
odata    equ     1280    ;location of output data for IFFT
locate   equ     $9C4    ;location of message to be sent (hex)

    sincos   points,coef
    opt      mex
```

```
      org      p:start
      move     #locate,r3
begin
      move     #data,r4
      do       #201,cod
      move     #output,r5
      jsr      readbit
      move     r4,x0
      move     #data+100,a
      cmp      x0,a       ;check to see if 100 symbols have been read
      jseq     outzero    ;output 55 zeros, read remaining symbols
      jsr      outbit
code
      org      p:startifft
      ifft     points,data,odata,coef   ;macro call to perform IFFT
;     jmp      begin

      nop
      swi

; The subroutine readbit reads the input value into memory one bit at a
; time

readbit
      move     #input+3,r2
      move                       y:(r3)+,a   ;read input symbol
      move     #>$1,x0
      do       #4,loop            ;symbol is read in one bit at a time
      and      x0,a        a,x1
      move     a1,x:(r2)-
      move     x1,a
      asr      a
loop
      rts

; The subroutine outbit outputs the values stored by readbit.  This four
; bit symbol is used as an index into a lookup table.

outbit
      move     #input,r2
      clr      b
      clr      a
      addl     b,a
      do       #4,loop2           ;symbol read out one bit at a time
      move     x:(r2)+,b0
      addl     b,a
loop2
      move     a0,r5              ;symbol is used to index into lookup table
      nop
      move     x:(r5),x0                   ;get real constellation value
      move                 y:(r5),y0   ;get imaginary constellation value
```

```
        move    x0,x:(r4)                       ;load real value into IFFT array
        move                    y0,y:(r4)+  ;load imag value into IFFT array nop
        rts


outzero
        do      #55,endzero                     ;subroutine loads zeros in center
        move    #0,x0                            ;55 locations of IFFT array to
        move    x0,x:(r4)                        ;allow for filtering
        nop
        move                    x0,y:(r4)+
endzero
        nop
        rts
```

APPENDIX F.   16 QAM BOUNDARY FILE

The following data file BOUN16.d contains the symbol data for each bounded area used in the 16 QAM decoder. It must be loaded into the DSP prior to running 16QAMDEC.

Data file 'BOUN16.D'

```
HX
00000200,0000020F
00000C,000008,000000,000004
00000D,00000A,000001,000006
00000E,000009,000002,000005
00000F,00000B,000003,000007
```

APPENDIX G.   16 QAM DECODER

```
; This program is a decoder for 16 QAM system.  After computing the FFT
; of the received data the points are read out of the array one at a time
; and compared with boundaries on the constellation to determine the
; received point.  The decoded symbols are then stored in memory.  It is
; assumed that the sampled input data is loaded into the input data
; location (1280) of the FFT.  If the sampled data is stored elsewhere a
; routine to load the FFT array will need to be added.
;
; The following registers are modified:
;
; r0      r4
; r1      r5
; r2      r6
;
; The following Data ALU registers are modified:
;
; a       b
; x0      x1
; y0      y1
;

qam16dec

      ident   1,1
      page 132,66,3,3,0
      opt nomd,nomex,loc,nocex,mu,cex

      include 'sincos'    ;include macro for sine cosine values
      include 'fft'       ;include macro for calculating FFT

      org 1:$0000

location dsm 16           ;storage locations for look up table
input    dsm 16           ;storage locations for input data

endlong   equ      *
      org x:endlong
storr6    ds 1
      org x:512
boundary1    ds    4      ;reserves four locations for boundary pts.
boundary2    ds    4      ;reserves four locations for boundary pts.
boundary3    ds    4      ;reserves four locations for boundary pts.
boundary4    ds    4      ;reserves four locations for boundary pts.

startfft equ   $100       ;starting location for fft program
```

```
start     equ   $250      ;starting location for decoder program
points    equ   256       ;number of points for FFT
coef      equ   1024      ;location of coefficients for FFT
data      equ   1280      ;location of input data for FFT
odata     equ   1536      ;location of output data for FFT


;Full scale values for 16 QAM constellation.

four     equ      $200000
three    equ      $180000
two      equ      $100000
one      equ      $080000
zero     equ      $000000
mone     equ      $F80000
mtwo     equ      $F00000
mthree   equ      $e80000
mfour    equ      $e00000

     sincos  points,coef
     opt   mex

     org   p:startfft
     fft   points,data,odata,coef      ;Macro call for FFT

     opt   cex
     org   p:start
     jsr   initialize
     do    #201,_endrun
     jsr   readdata
     jsr   findbound
_endrun
        nop
        swi


;This initialization routine initializes register and modifiers
;as well as clearing the memeory.
;The constellation is also loaded into memory here.

initialize
     move      #$ffff,m0 ;reset register to linear addressing
     move      #$ffff,m1 ;reset register to linear addressing
     move      #$ffff,m2 ;reset register to linear addressing
     move      #$ffff,m4 ;reset register to linear addressing
     move      #$ffff,m5 ;reset register to linear addressing
     move      #15,m6    ;set register for modulo 15 addressing
     move      #0,r1
     clr       b    #$0,r0
     clr       a    r0,r5
     do #50,clrmem
     move      a,x:(r0)+ b,y:(r5)+    ;clear first 50 memory locations
clrmem
```

```
        move    #input,b1
        move    b1,x:storr6

;Now load full scale values of the constellation in the table locations.


        move    #location,r0                    ;Real      Imag
        move    r0,r4
        move    #mone,a
        move    #mone,b
        move    a,x:(r0)+   b,y:(r4)+          ;-1        -1
        move    #mthree,b
        move    b,x:(r0)+   a,y:(r4)+          ;-3        -1
        move    a,x:(r0)+   b,y:(r4)+          ;-1        -3
        move    #mthree,a
        move    a,x:(r0)+   b,y:(r4)+          ;-3        -3
        move    #mone,b
        move    #one,a
        move    a,x:(r0)+   b,y:(r4)+          ; 1        -1
        move    #mthree,b
        move    a,x:(r0)+   b,y:(r4)+          ; 1        -3
        move    #three,b
        move    #mone,a
        move    b,x:(r0)+   a,y:(r4)+          ; 3        -1
        move    #mthree,a
        move    b,x:(r0)+   a,y:(r4)+          ; 3        -3
        move    #mthree,a
        move    #three,b
        move    #mone,x0
        move    #one,y1
        move    x0,x:(r0)+   y1,y:(r4)+        ;-1         1
        move    x0,x:(r0)+   b,y:(r4)+         ;-1         3
        move    a,x:(r0)+   y1,y:(r4)+         ;-3         1
        move    a,x:(r0)+   b,y:(r4)+          ;-3         3
        move    #one,a
        move    #one,x0
        move    #three,y1
        move    x0,x:(r0)+   a,y:(r4)+         ; 1         1
        move    b,x:(r0)+   a,y:(r4)+          ; 3         1
        move    x0,x:(r0)+   y1,y:(r4)+        ; 1         3
        move    b,x:(r0)+   y1,y:(r4)+         ; 3         3
        move    #odata,x0
        move    #$eff,y0
        move    x0,y:$424              ;store location for input to decoder
        move    y0,y:$425             ;store location for output of decoder
        rts

; The subroutine readdata reads in the data from the output of the FFT.
; The values are read out one point at a time.  The values are compared to
; boundaries on a partioned constellation to determine the received point.
```

```
readdata
      move      y:$424,r0
      nop
      move      r0,y0
      move      #odata+100,a      ;check to see if first 100 values have been
read
      cmp       y0,a
      jseq      _delzero          ;delete 55 zeros
      move      x:storr6,r6
      move      x:(r0),a          ;input data is loaded into storage location
      move      y:(r0)+,b         ;for decoding
      move      a,x:(r6)
      move      b,y:(r6)+
      move      r6,x:storr6
      move      r0,y:$424
      rts
;
;subroutine _delzero removes the zeros installed in the encoder
;
_delzero
      move      #odata+155,r0
      nop
      rts
;
; The subroutine findbound compares the value that has been read out of
; the array to the boundries on the constellation to decode the point.
; First the magnitude alone is used to find the correct bounded area then
; the signed values are used to determine the correct quadrant is used to
; increment the boundary pointer to find the correct point.
;

findbound

      move      x:-(r6),a       ; real value is stored in a
      move      #boundary1,r2   ; load starting position for bounded values
      move      #two,y0
      cmpm      y0,a            y:(r6),b   ; compare mag of real value to two
                                           ; imaginary value is stored in b
      jgt       bigtwo          ; x>2
      cmpm      y0,b            x:(r2),x0  ; compare mag of imag value to two
      jlt       continue                   ; x<2,y<2, load r2 with boundary 1 and
                                           ; continue
      move      #four,x1
      cmpm      x1,b            #boundary3,r2 ; compare magnitude of imag value
      nop
      move      x:(r2),x0
      jmp       continue                   ;x<2,y>2 and y<4, load r2 with boundary4

bigtwo
      cmpm      y0,b            #boundary2,r2 ;x>2 y<2
      nop
```

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```
        move      x:(r2),x0
        jlt       continue                ;x>2 and x<4, y<2 load boundary2 and
continue
        cmpm      x1,b          #boundary4,r2
        nop
        move      x:(r2),x0
        jmp       continue                ;x>2,y<4 load boundary 4 and continue

; This part of the routine finds the correct quadrant and updates the
; pointer to the correct point.

continue
        clr  a                 x:(r6),x1
        cmp  x1,a              y:(r6),y1
        jgt  negx
        cmp  y1,a              #3,n2
        jgt  posxnegy
posxposy
        jmp outputdata                ;output is in first quadrant
posxnegy
        move      x:(r2)+n2,x0    ;update r2 by 3, fourth quadrant
        jmp  outputdata
negx
        cmp  y1,a              #1,n2
        jgt  negxnegy
negxposy
        move      x:(r2)+n2,x0   ;update r2 by 1, second quadrant
        jmp  outputdata
negxnegy
        move #2,n2
        nop
        move x:(r2)+n2,x0         ;update r2 by 2, third quadrant

outputdata
        move      y:$425,r0
        move      x:(r2),a
        nop
        move      a1,y:(r0)+     ;move decoded symbol to output location
        move      r0,y:$425
        rts
```

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APPENDIX H. 2/3 RATE CODE DATA FILES

The following two data files are prescaled (full scale constellation values divided by 256) and conjugated. Both data files must be loaded into the DSP prior to running 23ENCOD. They are in the proper format to be loaded into memory.

Data file '23REAL.DAT'

```
HX
00000000,00000007
001000,FFF000,000800,FFF800
FFF800,000800,FFF000,001000
```

Data file '23IMAG.DAT'

```
HY
00000000,00000007
FFF800,000800,FFF000,001000
FFF000,001000,FFF800,000800
```

APPENDIX I. 2/3 RATE CODE ENCODER

```
; This program is an encoder for a 2/3 rate convolutional code.  A symbol
; consisting of two bits is read into the encoder, combinational
; logic is used to generate the third bit then a look up table is used
; to output the values from a point on the constellation.  The data files
; 23REAL.DAT and 23IMAG.DAT must first be loaded in memory prior to
; 23ENCOD.  The 256 values are loaded into an array and an IFFT is
; performed to generate the MFM signal.
;
; The following registers are modified;
;    r0      r3
;    r1      r4
;    r2      r5
;
; The following Data ALU registers are modified;
;  a   a0   a1
;  b   b0
;  x0  x1
;  y0  y1
;

23encod
        ident   1,1
     page   132,54
     opt     nomd,nomex,loc

     include  'sincos'
     include  'ifft'

start    equ       $100   ;starting program location
startifft    equ  $350   ;starting program location for IFFT
points   equ     256      ;number of points for IFFT
coef     equ     128      ;location of coefficients for IFFT
data     equ     768      ;location of input data for IFFT
odata    equ     1280     ;location of output data for IFFT
locate   equ     $9C4     ;starting location for message to be sent
input    equ     40       ;storage locations for input bits
output   equ     50       ;storage locations for encoded bits
statemen     equ   60   ;storage for past bits

     sincos  points,coef
     opt     mex

     org     p:start
     move    #locate,r3
```

```
begin
     move      #data,r4
      do       #201,code
     move      #output,r5
     jsr       readbit
     jsr       encode
     move      r4,x0
     move      #data+100,a
     cmp       x0,a        ;check to see if 100 symbols have been read
     jseq      outzero     ;output 55 zeros, read remaining symbols
     jsr       outbit
code
     org       p:startifft
     ifft      points,data,odata,coef;macro call to perform IFFT
;    jmp       begin

     nop
     swi

;
; The subroutine readbit reads two bits from the message and stores them
; in memory.
;
readbit
     move      #input+1,r2
     move                          y:(r3)+,a
     move      #>$1,x0
     do        #2,loop
     and       x0,a           a,x1
     move      a1,x:(r2)-
     move      x1,a
     asr       a
loop
     rts

;
; The subroutine encode performs the convolutional encoding to generate
; a redundant bit.
;
encode
     move      #input,r0
     move      #output,r5
     move      #statemem,r1
     move      x:(r0)+,x0          ;read bit x0
     move      x:(r0)-,x1          ;read bit x1
     move      x:(r1)+,a           ;read past bit s1
     move      x:(r1)-,b           ;read past bit s2
     eor       x1,a                ;x1 eor s1
     move      a,y1
     eor       y1,b                ;x1 eor s1 eor s2 = y2
     move      x0,y:(r5)           ;store x0
     move      x:(r1)+,a           ;read past bit s1
     move      x:(r1)-,x0          ;read past bit s2
```

```
        eor       x1,a                    ;x1 eor s1 = y1
        move      x0,x:(r1)+              ;move s2 to s1
        move      x1,x:(r1)              ;move x1 to s1
        move      a,x:(r0)+              ;temp store y1
        move      b,x:(r0)              ;temp store y2
        rts
```

; The subroutine outbit reads the three bits out of memory to form a
; symbol. The symbols are reconstructed by reading the bits one at a
; time added to to an empty register and shifted left until three bit
; symbol is formedA lookup table is used to get the values for the point
; cn the constellation.

```
outbit     move      #input,r2
        clr   b
        clr   a          y:(r5),b0
        addl      b,a
        do        #2,loop2
        move      x:(r2)+,b0
        addl      b,a
loop2
        move                          a0,r5        ;move symbol to r5
        nop
;
;use value in r5 to get points off the constellation
;
        move      x:(r5),x0
        move                          y:(r5),y0
;
;move constellation points to IFFT array
;
        move      x0,x:(r4)
        move                          y0,y:(r4)+
        nop
        rts
;
;subroutine outzero puts 55 zeros in IFFT array for filtering
;
outzero
        do        #55,endzero
        move      #0,x0
        move      x0,x:(r4)
        nop
        move                          x0,y:(r4)+
endzero
        nop
        rts
```

APPENDIX J. 2/3 RATE CODE BOUNDARY FILE

The following data file BOUN23.D contains the points on the constellation
used to find the minimum distances to states after the proper quadrant is
found.

Data file 'BOUN23.D

```
HX
00000200,0000020F
000040,000042,000044,000047
000041,000042,000044,000046
000041,000043,000045,000046
000040,000043,000045,000047
```

APPENDIX K.   2/3 RATE CODE VITERBI DECODER

```
; This program is a Viterbi Decoder for a 2/3 rate  convolutional encoder.
; There is a 16 time period delay which will approach the maximum possible
; gain for this type of encoder.  There are 64 locations needed in memory
; (16 past time periods x 4 states = 64).
;
; The following registers are modified:
;   r0    r5
;   r1    r6
;   r2    r7
;   r4
;
; The following Data ALU registers are modified:
;   a     b     x0    y0
;   a1    b1    x1    y1
;
; The following register modifiers are used;
;   m1    n0
;   m5    n2
;   m6
;
mfm23dec
      ident    1,1
      page 132,66,3,3,0
      opt nomd,nomex,loc,nocex,mu,cex

      include 'sincos'
      include 'fft'

      org 1:$0000

period    dsm     64          ;64 storage locations
location  dsm     8           ;constellation points
input     dsm     16          ;past 16 input bits
tables    dsm     4           ;accumlated distance
temp      dsm     4           ;temp storage for distance table

endlong   equ        *
      org x:endlong
storr6    ds      1
ynow      dsm     3           ;input bits
      org y:endlong
ypast     dsm     2           ;past bits
      org x:512
boundary1     ds   16         ;storage for boundary data
startfft      equ  $100       ;starting location for FFT
```

```
points          equ  256        ;number of points for FFT
coef            equ  1024       ;location of coefficients for FFT
data            equ  1280       ;location of sampled data for FFT
odata           equ  1536       ;output location for FFT
start           equ  $250       ;starting location for decoder program

; Define full scale constellation values.

two        equ     $100000
one        equ     $080000
zero       equ     $000000
mone       equ     $F80000
mtwo       equ     $F00000


large      equ  .9
small      equ  .1
offset     equ     $000000      ;can be used to distort data

     sincos  points,coef
     opt  mex

     org  p:startfft
     fft  points,data,odata,coef
     opt  cex
     org  p:start
     jsr  initialize
     do   #201,_endrun
     jsr  readdata
     jsr  findmindist
     jsr  accumdist
     jsr  traceback
     jsr  outputdata
_endrun
         nop
         swi
;
;this initialization routine initializes register and modifiers
;as well as clearing the memeory. The constellation is also loaded
;into memory here.  The accumulated distance array is set so that
;state zero starts out at a value of zero and all others start out
;larger, forcing the paths to merge at the zero states.
;
initialize
     move     #$ffff,m0       ;sets linear addressing
     move     #63,m1          ;sets modulo 63 addressing
     move     #$ffff,m2       ;sets linear addressing
     move     #$ffff,m4       ;sets linear addressing
     move     #63,m5          ;sets modulo 63 addressing
     move     #15,m6          ;sets modulo 16 addressing
     move     #0,r1
```

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```
        clr     b       #$0,r0
        clr     a       r0,r5
        do #256,clrmem
        move    a,x:(r0)+ b,y:(r5)+
clrmem
        move    #tables+1,r7
        move    #$400000,a1
        rep #3
        move    a1,x:(r7)+

        move    #input,b1
        move    b1,x:storr6

        move    #odata-16,r0
        move    r0,r2
        move    #0,x0
        move    #0,y0
        do      #16,_clrreg
        move    x0,x:(r0)+
        move    y0,y:(r2)+
_clrreg
;
;Now load full scale values of the constellation in the table locations.
;

        move    #location,r0            ;Real       Imag
        move    r0,r4
        move    #two,a
        move    #one,y1
        move    #mtwo,b
        move    #mone,y0
        move    #one,x1
        move    #mone,x0
        move    a,x:(r0)+   y1,y:(r4)+    ; 2          1
        move    x1,x:(r0)+  a,y:(r4)+     ;-2         -1
        move    x0,x:(r0)+  a,y:(r4)+     ; 1          2
        move    b,x:(r0)+   y1,y:(r4)+    ;-1         -2
        move    b,x:(r0)+   y0,y:(r4)+    ;-1          2
        move    x1,x:(r0)+  b,y:(r4)+     ; 1         -2
        move    x0,x:(r0)+  b,y:(r4)+     ; 1         -2
        move    a,x:(r0)+   y0,y:(r4)+    ; 2         -1
        move    #odata,x0
        move    #$eff,y0
        move    x0,y:$424
        move    y0,y:$425
        rts
;
; readdata reads in the data from the outpu of the FFT.  The data is read
; in as complex points on the constellation.
;
```

```
readdata
      move      y:$424,r0
      nop
      move      r0,y0
      move      #odata+100,a
      cmp       y0,a
      jseq      _delzero
      move      x:storr6,r6
      move      #>offset,x0
      move      x:(r0),a
      add       x0,a        y:(r0)+,b
      add       x0,b        a,x:(r6)
      move      b,y:(r6)+
      move      r6,x:storr6
      move      r0,y:$424
      rts


_delzero
      move      #odata+155,r0
      nop
      rts
;
; the minimum distance is found to the closest point in every state and
; stored. The values are stored so that indexing is made easier, state
; 0,2,3,1. This will greatly reduce the number of cycles needed later.
; a smoothing function is used to accumulate distances in the accumulated
; table so this minimum distance is multiplied by .1.
;
; The subroutine findmindist finds the quadrant of the received point
; which is used as a pointer into the boundary table to read the four
; closest points, one in each state.
;
findmindist
      move      #boundary1,r2
      clr   a   x:-(r6),x1
      cmp   x1,a          y:(r6),y1
      jgt   negx                        ;x<0
      cmp   y1,a          #12,n2
      jgt   posxnegy
posxposy
      jmp   findist
posxnegy
      move x:(r2)+n2,x0
      jmp   findist
negx
      cmp   y1,a          #4,n2
      jgt   negxnegy
negxposy
      move x:(r2)+n2,x0
      jmp   findist
```

```
negxnegy
     move #8,n2
     nop
     move x:(r2)+n2,x0
;
; findist finds the distance from the received point to the points read
; from the boundary table.
;
findist
     move      x:(r2)+,r0
     move      #tables,r4
     move      x:(r0),a
     sub       x1,a         y:(r0),b
     sub       y1,b         a,x0
     mpy       x0,x0,a      b,y0
     mac       y0,y0,a      x:(r2)+,r0
     move      #small,x0    a,y0
     mpy       x0,y0,a
     move      x:(r0),a     a,y:(r4)+
     sub       x1,a         y:(r0),b
     sub       y1,b         a,x0    y:(r4)+,y0
     mpy       x0,x0,a      b,y0
     mac       y0,y0,a      x:(r2)+,r0
     move      #small,x0 a,y0
     mpy       x0,y0,a
     move      x:(r0),a     a,y:(r4)+
     sub       x1,a         y:(r0),b
     sub       y1,b         a,x0
     mpy       x0,x0,a      b,y0
     mac       y0,y0,a      x:(r2)+,r0
     move      #small,x0 a,y0
     mpy       x0,y0,a
     move      x:(r0),a     a,y:(r4)-
     sub       x1,a         y:(r0),b
     sub       y1,b         a,x0y:(r4)-,y0
     mpy       x0,x0,a      b,y0
     mac       y0,y0,a      x:(r2)+,r0
     move      #small,x0 a,y0
     mpy       x0,y0,a
     move                   a,y:(r4)
     rts
;
;the accumulted distance routine  adds the smallest distance from the
;previously computed table for all pathes going into a state and
;does this for all four states.  Since only certain transitions are
;allowed the calculations are done in a specific order to reduce delay.
;
accumdist
     clr  a   #tables,r0
     move     #$7fffff,a1
     move     r0,r4
```

```
            move    #temp,r2
            move    #1,m0
            move    m0,m4
            move    #1,n1
            move    n1,n5
            move    r1,r5


        ;find minimum distance to state zero
            do    #2,statezero
            move x:(r0),x0     y:(r4),b
            add  x0,b
            cmp  b,a
            tge  b,a          r0,r3
            tge  b,a          r4,r7
            move x:(r0)+,x0    y:(r4)+,b
        statezero
            move r3,x:(r1)+n1
            move a,x:(r2)+     y:(r4)+,b
            clr  a            r7,y:(r5)+n5
            move #$7fffff,a1


        ;find minimum distance to state two
            do    #2,statetwo
            move x:(r0),x0     y:(r4),b
            add  x0,b
            cmp  b,a
            tge  b,a          r0,r3
            tge  b,a          r4,r7
            move x:(r0)+,x0    y:(r4)+,b
        statetwo
            move r3,x:(r1)+n1
            move a,x:(r2)+     y:(r4)+,b
            clr  a            r7,y:(r5)+n5
            move             #tables+2,r4
            move             r4,r0
            move             x:(r1)-n1,a
            clr  a           x:(r1)-,b
            move #$7fffff,a1
            move r1,r5

        ;find minimum distance to state one
            do    #2,stateone
            move x:(r0),x0     y:(r4),b
            add  x0,b
            cmp  b,a
            tge  b,a     r0,r3
            tge  b,a     r4,r7
```

```
        move x:(r0)+,x0    y:(r4)-,b
stateone
        move r3,x:(r1)+n1
        move a,x:(r2)+
        clr  a             r7,y:(r5)+n5
        move #$7fffff,a1

;find minimum distance to state three
        do   #2,statethree
        move x:(r0),x0     y:(r4),b
        add  x0,b
        cmp  b,a
        tge  b,a      r0,r3
        tge  b,a      r4,r7
        move x:(r0)+,x0    y:(r4)-,b
statethree
        move r3,x:(r1)+n1
        move a,x:(r2)+
        clr  a             r7,y:(r5)+n5
        move #$7fffff,a1
        move (r4)+

;now move new accumulated distances into the  accumulated distance
;table from the temporary table
;also find the min distance state and store in r4 which is no longer used

        move      #$ffff,m0
        move      #$ffff,m4
        move      #temp,r3
        move      #tables,r0
        move      #large,x1
        move      #2,n0
        do        #3,endtable
        move      x:(r3)+,x0
        mpy       x1,x0,a
        cmp  a,b            a,x:(r0)+n0
        tge  a,b            r0,r4
endtable

        move      #tables+1,r0
        do        #4,endtablex
        move      x:(r3)+,x0
        mpy       x1,x0,a
        cmp       a,b       a,x:(r0)+n0
        tge       a,b       r0,r4
endtablex

;store in r0 instead of r4

        move      r4,r0
        move      #4,n1
```

```
        move    (r0)-n0
        rts


;the traceback routine now goes back through every time period starting
;with the current time period and finds the state from which the path
;came from one time period previous.  At the end of this search, the
;last state found will also point to the path at that state, which is the
;output  of the trellis.

traceback

;find the displacement from the pointer to table and store value in n4

        move    #tables,n0
        move    (r1)-n1
        lua     (r0)-n0,n5
        move    r1,r5
        do      #15,endtrace
        move    (r1)-n1
        move    x:(r5+n5),r0
        move    r1,r5
        lua     (r0)-n0,n5
endtrace
        move    #location,r0
        move    y:(r5+n5),a
        rts


;the output data routine unscrambles the path order and finds one
;of the two points on the constellation coresponding to the output state
;which is closest to the original input at that time period.

outputdata
        move  a,b
        move  #>$b1,x0
        cmp  x0,a      #>$b2,y0
        teq  y0,b
        cmp  y0,a      #>$b3,x0
        teq  x0,b
        cmp  x0,a      #>$b1,y0
        teq  y0,b
        move #>$b5,x0
        cmp  x0,a      #>$b7,y0
        teq  y0,b
        cmp  y0,a
        teq  x0,b
        move b,r2
        move #tables,n2
        move x:storr6,r6
        lua  (r2)-n2,n3
        move n3,a
```

```
        asl   a
        asl   a
        move  a,n0
        move  r6,r3
        lua   (r0)+n0,r4
        move  #>$7fffff,x1
        move  r4,r0
        do    #4,endout
        move  x:(r3),a      y:(r6),b
        move  x:(r0)+,x0    y:(r4)+,y0
        sub   x0,a
        sub   y0,b          a,x0
        mpy   x0,x0,a       b,y0
        mac   y0,y0,a
        tfr   a,b           x1,a
        cmp   x1,b
        tlt   b,a           r0,r7
        move  a,x1
endout
        clr   a   (r7)-
        move      #location,n0
        move      r7,r0
        move      #$f,a1
        lua       (r0)-n0,r7
        move      r7,x0
        and       x0,a
        jsr       convoldec

        move      y:$425,r0
        nop
        move      a0,y:(r0)+
        move      r0,y:$425
        rts
```

;The subroutine convoldec decodes the received symbol by using
;combinational logic.

```
convoldec
        move      #ynow+2,r0
        move      #>$1,x0
        move      #ypast,r7
        do        #3,loop
        and       x0,a        a,x1
        move      a1,x:(r0)-
        move      x1,a
        asr       a
loop

        move x:(r0)+,a      y:(r7)+,y0      ;read s1
        move x:(r0)+,b      y:(r7)-,y1      ;read y0 and s2
        move x:(r0)-,a                      ;read y1
```

60

```
      nop
      eor  y0,a                               ;yl eor sl =xl
      move                   yl,y:(r7)+       ;update past states
      move a,y:(r7)-
      move b,x:(r0)-
      move a,x:(r0)+
      clr  a
      clr  b
      do   #2,loop2
      move x:(r0)-,b0                         ;output decoded bits
      addl b,a
loop2
      rts
```

APPENDIX L.   V.32 DATA FILES

The following data files are prescaled (full scale constellation values divided by 256) and conjugated. Both data files must be loaded into the DSP prior to running MFMENC. They are properly formatted to be read into memory.

Data file 'QAMREAL.DAT'

```
HX
00000000,0000001F
FFE000,000000,000000,002000
002000,000000,000000,FFE000
FFF000,FFF000,001000,001000
001000,001000,FFF000,FFF000
FFE800,000800,FFE800,000800
001800,FFF800,001800,FFF800
000800,FFE800,000800,000800
FFF800,001800,FFF800,FFF800
```

Data file 'QAMIMAG.DAT'

```
HY
00000000,0000001F
FFF800,001800,FFF800,FFF800
000800,FFE800,000800,000800
FFE800,000800,FFE800,000800
001800,FFF800,001800,FFF800
001000,001000,FFF000,FFF000
FFF000,FFF000,001000,001000
FFE000,000000,000000,002000
002000,000000,00C000,FFE000
```

APPENDIX M. V.32 ENCODER

```
; This is a an encoder for V.32 standard with differential encoding.
; The encoder can also be used with the differential removed by commenting
; out the call to diff.

mfmencod
        ident   1,1
      page    132,54
      opt     nomd,nomex,loc

      include  'sincos'
      include  'ifft'


        org   x:$40
statemem      ds    3               ;set 3 locations for p st states
input         ds    4               ;set up 4 locations for input bits


        org   y:$40
ylpast        ds    2               ;set up two locations for diff encoder
output        ds    1               ;output of convolutional encoder

start       equ     $100            ;starting location for encoder
startifft   equ     $350            ;starting location for IFFT
points   equ    256                 ;number of points for IFFT
odata    equ    1280                ;output of IFFT
data     equ    768                 ;input data for IFFT
coef     equ    128                 ;location for coeficients for IFFT
locate   equ    $9C4                ;starting location for message


      sincos   points,coef
      opt      mex

        org      p:start
      move       #ylpast,r5
      move       #statemem,r3
      move        #locate,r6

begin
      move       #data,r7
      do         #201,code          ;reads 201 message symbols
      move       #output,r4
      jsr        readbit
      jsr        diff
      jsr        encode
```

```
        move    r7,x0
        move    #data+100,a
        cmp     x0,a                        ;checks if first 100 symbols read
        jseq    outzero                     ;outputs 55 zeros
        jsr     outbit
code
        org     p:startifft
        ifft    points,data,odata,coef
;       jmp     begin

        nop
        swi


;
;subroutine readbit reads in the four bit symbol one bit at a time
;
readbit
        move #input+3,r2
        move                y:(r6)+,a
        move #>$1,x0
        do   #4,loop
        and  x0,a           a,x1
        move al,x:(r2)-
        move x1,a
        asr  a
loop
        rts
;
;subroutine diff differentially encodes the two most significant bits
;

diff
        move    #input,r1
        move                y:(r5)+,y0
        move x:(r1)+,a      y:(r5)-,y1
        move x:(r1)-,b
        eor  y0,a           a,x0
        eor  y1,b           a,x1
        move                x0,a
        and  y0,a           b,y1
        eor  y1,a           x1,b
        move b,x:(r1)+      b,y:(r5)+
        move a,x:(r1)       a,y:(r5)-
        rts


;
;subroutine encode convolutionally encodes the four bits to generate a
;fifth bit
;
encode
        move    #input,r0
```

```
        move     #output,r4
        move     #statemem,r1
        move     x:(r0)+,x1
        move     x:(r1)+,a
        move                    a,y:(r4)
        and  x1,a               x:(r0),x0
    move x:(r1)-,b
        eor  x0,b               a,y0
        eor  y0,b               b,y1
        move b,x:(r1)+          y:(r4),b
        and  y1,b               x0,a
        move (r1)+
        eor  x1,a               x:(r1),x0
        eor  x0,a               y:(r4),y1
        move b,y0
        eor  y0,a               y1,x:(r1)-
        move a,x:(r1)+
        rts
;
;subroutine outbit reads the five bit symbol and uses it to index into the
;lookup table to get the values of the point on the constellation
;
outbit     move      #input,r2
        clr  b
        clr  a                  y:(r4),b0
        addl b,a
        do   #4,loop2
        move x:(r2)+,b0
        addl b,a
loop2
        move                    a0,r4
        nop
        move x:(r4),x0
        move                    y:(r4),y0
;
;move the values read from the constellation to the input array for the
;IFFT
;
    move x0,x:(r7)
        move                    y0,y:(r7)+
    nop
    rts
;
;subroutine outzero loads 55 zeros into the IFFT array for filtering
;
outzero
        do   #55,endzero
        move #0,x0
    move x0,x:(r7)
    nop
```

```
        move              x0,y:(r7)+
endzero
        nop
        rts
```

APPENDIX N.   V.32 BOUNDARY FILE

This data file BOUND.D contains the eight closest points (one in each state) to each of the 52 bounded areas used in the partitioned constellation. It must be loaded prior to using MFMDIFDE.

Data file 'BOUND.D'

```
HX
00000200,0000039F
000082,000086,00008B,00008D
000093,G00095,00009A,00009E
000082,000086,000089,00008F
000093,000095,00009A,00009E
000082,000086,000089,C0008F
000091,000097,00009A,00009E
000082,000086,00008B,00008D
000091,000097,00009A,00009E
000082,000086,00008B,00008D
000093,000094,00009A,00009D
000082,000086,000089,00008F
000092,000095,000099,00009E
000082,000086,000089,00008F
000090,000097,000099,00009E
000082,000086,00008B,00008D
000091,000096,00009A,00009D
000083,000084,00008B,00008D
000093,000094,00009A,00009D
000080,000087,000089,00008F
000092,000095,000099,00009E
000080,000087,000089,00008F
000090,000097,000099,00009E
000083,000084,00008B,00008D
000091,000096,00009A,00009D
000082,000085,00008A,00008D
000093,000095,00009A,00009E
000082,000085,000088,00008F
000093,000095,00009A,00009E
000081,000086,000089,00008E
000091,000097,00009A,00009E
000081,000086,00008B,00008C
000091,000097,00009A,00009E
000082,000085,00008A,00008D
000093,000094,00009A,00009D
000082,000085,000088,00008F
000092,000095,000099,00009E
000081,000086,000089,00008E
```

```
000090,000097,000099,00009E
000081,000086,00008B,00008C
000091,000096,00009A,00009D
000082,000085,00008A,00008D
000093,000095,000098,00009F
000082,000085,000088,00008F
000093,000095,000098,00009F
000081,000086,000089,00008E
000091,000097,00009B,00009C
000081,000086,00008B,00008C
000091,000097,00009B,00009C
000083,000084,00008A,00008D
000093,000094,00009A,00009D
000080,000087,000088,00008F
000092,000095,000099,00009E
000080,000087,000089,00008E
000090,000097,000099,00009E
000083,000084,00008B,00008C
000091,000096,00009A,00009D
000082,000085,00008A,00008D
000093,000094,000098,00009F
000082,000085,000088,00008F
000092,000095,000098,00009F
000081,000086,000089,00008E
000090,000097,00009B,00009C
000081,000086,00008B,00008C
000091,000096,00009B,00009C
000083,000085,00008A,00008D
000093,000094,000098,00009D
000080,000085,000088,00008F
000092,000095,000099,00009F
000081,000087,000089,00008E
000090,000097,000099,00009C
000081,000084,00008B,00008C
000091,000096,00009B,00009D
000082,000085,00008A,00008D
000093,000094,000098,00009D
000082,000085,000088,00008F
000092,000095,000099,00009F
000081,000086,000089,00008E
000090,000097,000099,00009C
000081,000086,00008B,00008C
000091,000096,00009B,00009D
000083,000085,00008A,00008D
000093,000094,00009A,00009D
000080,000085,000088,00008F
000092,000095,000099,00009E
000081,000087,000089,00008E
000090,000097,000099,00009E
000081,000084,00008B,00008C
000091,000096,00009A,00009D
```

```
000083,000084,00008A,00008D
000093,000094,000098,00009D
000080,000087,000088,00008F
000092,000095,000099,00009F
000080,000087,000089,00008E
000090,000097,000099,00009C
000083,000084,00008B,00008C
000091,000096,00009B,00009D
000083,000085,00008A,00008D
000093,000094,000098,00009F
000080,000085,000088,00008F
000092,000095,000098,00009F
000081,000087,000089,00008E
000090,000097,00009B,00009C
000081,000084,00008B,00008C
000091,000096,00009B,00009C
```

```
;This program is a Viterbi Decoder for V.32.  There is a 16
;time period delay which will approach the maximum possible
;gain for this type of encoder. If the differential encoder
;was not used in the encoder than the call to diff must be
;commented out.
;
;There are 128 memory locations allocated for path memory (8 states x 16
;time periods =128 locations).  The full scale constellation values are
;loaded into memory during the intialization routine.
;
mfmdecod
      ident   1,1
      page 132,66,3,3,0
      opt nomd,nomex,loc,nocex,mu,cex

      include 'sincos'
      include 'fft'

      org 1:$0000

period    dsm      128         ;128 locations for path memory
location dsm       32          ;32 locations for constellation points
input     dsm      16          ;16 locations for input points
tables    dsm      8           ;8 locations for acumulated distance table
temp      dsm      8           ;8 temp locations for distances
endlong   equ      *
      org x:endlong
storr6    ds 1
ynow      ds 4                 ;4 locations for input bits
      org y:endlong
ypast     ds 2                 ;2 past bits for differential decoder


;
;13 boundary tables with 8 points in each of the 4 quadrants
;
      org x:512
boundry1        ds   32
boundry2        ds   32
boundry3        ds   32
boundry4        ds   32
boundry5        ds   32
boundry6        ds   32
boundry7        ds   32
boundry8        ds   32
```

```
boundry9        ds   32
boundry10       ds   32
boundry11       ds   32
boundry12       ds   32
boundry13       ds   32


startfft        equ        $100        ;starting location for FFT routine
points    equ      256                 ;number of points for FFT
coef      equ      1024                ;location of FFT coefficients
data      equ      1280                ;location of sampled data
odata     equ      1536                ;output data from FFT
start     equ      $250                ;starting location of decoder
;
;Load in full scale constellation values
;
four      equ        $200000
three     equ        $180000
two       equ        $100000
one       equ        $080000
zero      equ        $000000
mone      equ        $F80000
mtwo      equ        $F00000
mthree   ·equ        $e80000
mfour     equ        $e00000

large     equ        .9
small     equ        .1
offset    equ        $000000

     sincos    points,coef
     opt  mex

     org  p:startfft
     fft  points,data,odata,coef
     opt  cex
     org  p:start
     jsr  initialize
     do   #217,_endrun
      jsr  readdata
      jsr  findmindist
      jsr  accumdist
     jsr  traceback
      jsr  outputdata
_endrun
        nop
        swi

;
; this initialization routine initializes register and
; modifiers as well as clearing the memeory. The constellation
; is also loaded into memory here. The accumulated distance
; array is set so that state zero starts out at a value of
```

```
; zero and all others start out larger, forcing the paths
; to merge at the zero states.
;
initialize
     move #$ffff,m0              ;linear addressing
     move #127,m1               ;modulo 127 addressing
     move #$ffff,m2             ;linear addressing
     move #$ffff,m4             ;linear addressing
     move #127,m5               ;modulo 127 addressing
     move #15,m6                ;modulo 15 addressing
     move #0,r1
     clr   b     #$0,r0
    clr   a     r0,r5
     do #256,clrmem
     move a,x:(r0)+     b,y:(r5)+
clrmem
     move     #tables+1,r7
     move     #$400000,a1
     rep #7
     move     a1,x:(r7)+

     move     #input,b1
     move     b1,x:storr6

     move     #odata-16,r0
     move     r0,r2
     move     #0,x0
     move     #0,y0
     do       #16,_clrreg
     move     x0,x:(r0)+
     move     y0,y:(r2)+
_clrreg

; Now load full scale values of the constellationin the table
; location.


     move     #location,r0
     move     r0,r4
     move     #mfour,a                           ;Real      Imag
     move     #one,b
     move     a,x:(r0)+     b,y:(r4)+            ; -4        1
     move     #zero,a
     move     #mthree,b
     move     a,x:(r0)+     b,y:(r4)+            ;  0        -3
     move     #one,b
     move     a,x:(r0)+     b,y:(r4)+            ;  0        1
     move     #four,a
     move     a,x:(r0)+     b,y:(r4)+            ;  4        1
     move     #mone,b
     move     a,x:(r0)+     b,y:(r4)+            ;  4        -1
```

```
move    #zero,a
move    #three,b
move    a,x:(r0)+    b,y:(r4)+           ;  0         3
move    #mone,b
move    a,x:(r0)+    b,y:(r4)+           ;  0        -1
move    #mfour,a
move    a,x:(r0)+    b,y:(r4)+           ; -4        -1
move    #mtwo,a
move    #three,b
move    #mone,y1
move    a,x:(r0)+    b,y:(r4)+           ; -2         3
move    a,x:(r0)+    y1,y:(r4)+          ; -2        -1
move    #two,a
move    a,x:(r0)+    b,y:(r4)+           ;  2         3
move    a,x:(r0)+    y1,y:(r4)+          ;  2        -1
move    #one,b
move    #mthree,y1
move    a,x:(r0)+    y1,y:(r4)+          ;  2        -3
move    a,x:(r0)+    b,y:(r4)+           ;  2         1
move    #mtwo,a
move    a,x:(r0)+    y1,y:(r4)+          ; -2        -3
move    a,x:(r0)+    b,y:(r4)+           ; -2         1
move    #one,a
move    a,x0
move    #mthree,a
move    #two,b
move    b,y0
move    #mtwo,b
move    a,x:(r0)+    b,y:(r4)+           ; -3        -2
move    x0,x:(r0)+   b,y:(r4)+           ;  1        -2
move    a,x:(r0)+    y0,y:(r4)           ; -3         2
move    x0,x:(r0)+   y0,y:(r4)+          ;  1         2
move    #three,a
move    a,x0
move    #mone,a
move    x0,x:(r0)+   y0,y:(r4)+          ;  3         2
move    a,x:(r0)+    y0,y:(r4)+          ; -1         2
move    x0,x:(r0)+   b,y:(r4)+           ;  3        -2
move    a,x:(r0)+    b,y:(r4)+           ; -2        -1
move    #one,a
move    #zero,b
move    b,y0
move    #four,b
move    a,x:(r0)+    b,y:(r4)+           ;  1         4
move    #mthree,x0
move    x0,x:(r0)+   y0,y:(r4)+          ; -3         0
move    a,x:(r0)+    y0,y:(r4)+          ;  1         0
move    #mfour,b
move    a,x:(r0)+    b,y:(r4)+           ;  1        -4
move    #mone,a
move    a,x:(r0)+    b,y:(r4)+           ; -1        -4
```

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```
      move    #three,x0
      move    x0,x:(r0)+    y0,y:(r4)+        ;  3         0
      move    a,x:(r0)+     y0,y:(r4)+        ; -1         0
      move    #four,b
      move    a,x:(r0)+     b,y:(r4)+         ; -1         4
      move    #odata,x0
      move    #$eff,y0
      move    x0,y:$424
      move    y0,y:$425
      rts
;
;readdata reads in the data from the output of the FFT.
;
readdata
      move y:$424,r0
      nop
      move r0,y0
      move #odata+100,a
      cmp  y0,a
      jseq _delzero
      move x:storr6,r6
      move #>offset,x0
      move x:(r0),a
      add  x0,a             y:(r0)+,b
      add  x0,b             a,x:(r6)
      move b,y:(r6)+
      move r6,x:storr6
      move r0,y:$424
      rts


_delzero
      move    #odata+155,r0
      nop
      rts
;
;the minimum distance is found to the closest point in every
; state and stored. The values are stored so that indexing is
; made easier, state 0,2,3,1,4,7,6,5. This will greatly reduce
; the number of cycles needed later. A smoothing function is
; used to accumulate distances in the accumulated table so
; this minimum distance is multiplied by .1.
;
;The subroutine findmindist compares the received points to boundaries
;on the constellation.  Once a bounded area is found the closest eight
;points are read out of the boundary data file and used to update the
;path distances.
;
findmindist

      move x:-(r6),a
      move #one,x0
```

```
        cmpm x0,a      y:(r6),b
        jgt  bigone                    ;x>1
        cmpm x0,b      #boundry1,r2
        jlt  continue                  ;x<1,y<1, load r2 with boundry 1
                                       ;and continue

        move #two,x1
        cmpm x1,b      #boundry4,r2
        jlt  continue                  ;x<1,y>1andy<2, load r2 with
                                       ;boundry4, go on

        move           #boundry6,r2
        jmp  continue                  ;x<1,y>2, load r2 with boundry6
                                       ;and continue
bigone    move #two,x1
        cmpm x1,a
        jgt  bigtwo                    ;x>2, jmp to that case
        cmpm x0,b      #boundry2,r2
        jlt  continue                  ;x>1 and x<2, y<1 load boundry2
                                       ;and continue

        cmpm x1,b      #boundry5,r2
        jlt  continue                  ;x>1,y<2 load boundry 5 and
                                       ;continue
bigtwo
        cmpm x0,b      #boundry3,r2
        jlt  continue                  ;x>2 and y<1 so load boundry3
                                       ;and continue

        abs  a         #two,y0
        abs  b         a,x1
        sub  y0,a      b,y1
        sub  x0,b
        cmpm a,b       y1,b
        jgt  greatery1
        cmp  y0,b      #boundry7,r2
        jlt  continue
        move           #boundry12,r2
        jmp  continue
greatery1
        sub  y0,b      x1,a
        sub  x0,a
        cmpm a,b       x1,a
        jgt  greatery2
        cmp  y0,a      #boundry10,r2
        jlt  continue
        move y1,b
        cmp  y0,b      #boundry11,r2
        jlt  continue
        move           #boundry9,r2
        jmp  continue
greatery2
        cmp  y0,a      #boundry8,r2
        jlt  continue
        move           #boundry13,r2
```

```
continue
      clr   a          x:(r6),x1
      cmp   x1,a       y:(r6),y1
      jgt   negx
      cmp   y1,a       #24,n2
      jgt   posxnegy
posxposy
      jmp  findist
posxnegy
      move x:(r2)+n2,x0                    ;update r2 by 24
      jmp  findist
negx
      cmp y1,a         #8,n2
      jgt negxnegy
negxposy
      move x:(r2)+n2,x0                    ;update r2 by 8
      jmp  findist
negxnegy
      move x:(r2)+n2,x0                    ;update r2 by 16
      move x:(r2)+n2,x0
;
;The subroutine findist finds the euclidean distance between the received
;point and the eight points read out of the boundary table.  The x and y
;coordinates are subtracted, squared and added.  The square root is not
;performed.
:

findist
      move x:(r2)+,r0
      move #tables,r4
      move x:(r0),a
      sub  x1,a        y:(r0),b
      sub  y1,b        a,x0
      mpy  x0,x0,a     b,y0
      mac  y0,y0,a     x:(r2)+,r0
      move #small,x0   a,y0
      mpy  x0,y0,a
      move x:(r0),a    a,y:(r4)+
      sub  x1,a        y:(r0),b
      sub  y1,b        a,x0     y:(r4)+,y0
      mpy  x0,x0,a     b,y0
      mac  y0,y0,a     x:(r2)+,r0
      move #small,x0   a,y0
      mpy  x0,y0,a     y:(r4)+,b
      move x:(r0),a    a,y:(r4)-
      sub  x1,a        y:(r0),b
      sub  y1,b        a,x0     y:(r4)-,y0
      mpy  x0,x0,a     b,y0
      mac  y0,y0,a     x:(r2)+,r0
      move #small,x0   a,y0
      mpy  x0,y0,a
```

```
        move x:(r0),a       a,y:(r4)+
        sub  x1,a           y:(r0),b
        sub  y1,b           a,x0
        mpy  x0,x0,a        b,y0
        mac  y0,y0,a        x:(r2)+,r0
        move #small,x0       a,y0
        mpy  x0,y0,a
        move x:(r0),a       a,y:(r4)+
        sub  x1,a           y:(r0),b
        sub  y1,b           a,x0   y:(r4)+,y0
        mpy  x0,x0,a        b,y0
        mac  y0,y0,a        x:(r2)+,r0
        move #small,x0       a,y0
        mpy  x0,y0,a
        move x:(r0),a       a,y:(r4)+
        sub  x1,a           y:(r0),b
        sub  y1,b           a,x0   y:(r4)+,y0
        mpy  x0,x0,a        b,y0
        mac  y0,y0,a        x:(r2)+,r0
        move #small,x0       a,y0
        mpy  x0,y0,a        y:(r4)+,b
        move x:(r0),a       a,y:(r4)-
        sub  x1,a           y:(r0),b
        sub  y1,b           a,x0
        mpy  x0,x0,a        b,y0
        mac  y0,y0,a        x:(r2)+,r0
        move #small,x0       a,y0
        mpy  x0,y0,a
        move x:(r0),a       a,y:(r4)-
        sub  x1,a           y:(r0),b
        sub  y1,b           a,x0
        mpy  x0,x0,a        b,y0
        mac  y0,y0,a        x:(r2)+,r0
        move #small,x0       a,y0
        mpy  x0,y0,a
        move a,y:(r4)
        rts

;the accumulted distance routine  adds the smallest distance
;from the previously computed table for all pathes going into
;a state and does this for all eight states.

accumdist
        clr  a    #tables,r0
        move      #$7fffff,a1
        move      r0,r4
        move      #temp,r2
        move      #3,m0
        move      m0,m4
        move      #2,n1
        move      n1,n5
```

```
        move    r1,r5

;Distances in the accumulated distance table are added to distances in the
;path table and compared for the four paths.  This is done by incrementing
;through a specially ordered path table.

;find minimum distance to state zero
        do      #4,statezero
        move x:(r0),x0      y:(r4),b
        add  x0,b
        cmp  b,a
        tge  b,a        r0,r3
        tge  b,a        r4,r7
        move x:(r0)+,x0     y:(r4)+,b
statezero
        move r3,x:(r1)+n1
        move a,x:(r2)+      y:(r4)+,b
        clr  a          r7,y:(r5)+n5
        move #$7fffff,a1




;find minimum distance to state two
        do      #4,statetwo
        move x:(r0),x0      y:(r4),b
        add  x0,b
        cmp  b,a
        tge  b,a        r0,r3
        tge  b,a        r4,r7
        move x:(r0)+,x0     y:(r4)-,b
statetwo
        move r3,x:(r1)+n1
        move a,x:(r2)+      y:(r4)+,b
        clr  a          r7,y:(r5)+n5
        move #$7fffff,a1

;find minimum distance to state four
        do      #4,statefour
        move x:(r0),x0      y:(r4),b
        add  x0,b
        cmp  b,a
        tge  b,a        r0,r3
        tge  b,a        r4,r7
        move x:(r0)+,x0     y:(r4)+,b
statefour
        move r3,x:(r1)+n1
        move a,x:(r2)+      y:(r4)+,b
        clr  a          r7,y:(r5)+n5
        move #$7fffff,a1

;find minimum distance to state six
```

```
        do    #4,statezsix
        move x:(r0),x0       y:(r4),b
        add   x0,b
        cmp   b,a
        tge   b,a       r0,r3
        tge   b,a       r4,r7
        move x:(r0)+,x0      y:(r4)-,b
statezsix
        move r3,x:(r1)-n1
        move a,x:(r2)+
        move r7,y:(r5)
        move #tables+4,r4
        move r4,r0
        move x:(r1)-n1,a
        clr  a    x:(r1)-,b
        move #$7fffff,a1
        move r1,r5

;find minimum distance to state one
        do    #4,stateone
        move x:(r0),x0       y:(r4),b
        add   x0,b
        cmp   b,a
        tge   b,a       r0,r3
        tge   b,a       r4,r7
        move x:(r0)+,x0      y:(r4)+,b
stateone
        move r3,x:(r1)+n1
        move a,x:(r2)+       y:(r4)+,b
        clr  a    r7,y:(r5)+n5
        move #$7fffff,a1

;find minimum distance to state three
        do    #4,statethree
        move x:(r0),x0       y:(r4),b
        add   x0,b
        cmp   b,a
        tge   b,a       r0,r3
        tge   b,a       r4,r7
        move x:(r0)+,x0      y:(r4)-,b
statethree
        move r3,x:(r1)+n1
        move a,x:(r2)+       y:(r4)+,b
        clr  a    r7,y:(r5)+n5
        move #$7fffff,a1
        move (r4)+

;find minimum distance to state five
        do    #4,statefive
        move x:(r0),x0       y:(r4),b
        add   x0,b
```

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```
      cmp  b,a
      tge  b,a        r0,r3
      tge  b,a        r4,r7
      move x:(r0)+,x0    y:(r4)-,b
statefive
      move r3,x:(r1)+n1
      move a,x:(r2)+     y:(r4)-,b
      clr  a   r7,y:(r5)+n5
      move #$7fffff,a1

;find minimum distance to state seven
      do   #4,stateseven
      move x:(r0),x0     y:(r4),b
      add  x0,b
      cmp  b,a
      tge  b,a      r0,r3
      tge  b,a      r4,r7
      move x:(r0)+,x0    y:(r4)+,b
stateseven
      move r3,x:(r1)+
      move a,x:(r2)+     y:(r4)+,b
      clr  b   r7,y:(r5)+
      move #$7fffff,b1

;now move new accumulated distances into the  accumulated
;distance table from the temporary table also find the min
;distance state and store in r4 which is no longer used

      move #$ffff,m0
      move #$ffff,m4
      move #temp,r3
      move #tables,r0
      move #large,x1
      move #2,n0
      do   #4,endtable
      move x:(r3)+,x0
      mpy  x1,x0,a
      cmp  a,b       a,x:(r0)+n0
      tge  a,b       r0,r4
endtable
      move #tables+1,r0
      do   #4,endtablex
      move x:(r3)+,x0
      mpy  x1,x0,a
      cmp  a,b       a,x:(r0)+n0
      tge  a,b       r0,r4
endtablex
;
;store in r0 instead of r4
;
      move    r4,r0
```

```
        move    #8,n1
        move    (r0)-n0
        rts


;the traceback routine now goes back through every time period
;starting with the current time period and finds the state
;from which the path came from one time period previous.  At
;the end of this search, the last state found will also point
;to the path at that state, which is the output  of the
;trellis.

traceback

;find the displacement from the pointer to table and store
;value in n4

        move #tables,n0
        move (r1)-n1
        lua  (r0)-n0,n5
        move r1,r5
        do   #15,endtrace
        move (r1)-n1
        move x:(r5+n5),r0
        move r1,r5
        lua  (r0)-n0,n5
endtrace
        move #location,r0
        move y:(r5+n5),a
        rts


;the output data routine unscrambles the path order and finds
;one of the four points on the constellation coresponding to
;the output state which is closest to the original input at
;that time period.

outputdata
        move a,b
        move            #>$b1,x0
        cmp  x0,a       #>$b2,y0
        teq  y0,b
        cmp  y0,a       #>$b3,x0
        teq  x0,b
        cmp  x0,a       #>$b1,y0
        teq  y0,b
        move            #>$b5,x0
        cmp  x0,a       #>$b7,y0
        teq  y0,b
        cmp  y0,a
        teq  x0,b
        move b,r2
```

```
        move #tables,n2
        move x:storr6,r6
        lua  (r2)-n2,n3
        move n3,a
        asl  a
        asl  a
        move a,n0
        move r6,r3
        lua  (r0)+n0,r4
        move #>$7fffff,x1
        move r4,r0
        do   #4,endout
        move x:(r3),a        y:(r6),b
        move x:(r0)+,x0      y:(r4)+,y0
        sub  x0,a
        sub  y0,b       a,x0
        mpy  x0,x0,a  b,y0
        mac  y0,y0,a
        tfr  a,b        x1,a
        cmp  x1,b
        tlt  b,a        r0,r7
        move a,x1
endout
        clr  a    (r7)-
        move #location,n0
        move r7,r0
        move #$f,a1
        lua  (r0)-n0,r7
        move r7,x0
        and  x0,a
        jsr  diff
        move y:$425,r0
        nop
        move a0,y:(r0)+
        move r0,y:$425
        rts
;
;The subroutine diff differentially decodes the two most significant bits.
;Each bit is stored in its own memory and the bits are decoded using
;Q1n = Y1n EOR Y1n-1, Q2n = (Q1n AND Y1n-1) EOR Y2n-1 EOR Y2n.  The four
;bit symbol is formed and output.
;
diff
        move #ynow+3,r0
        move #>$1,x0
        move #ypast,r7
        do   #4,diffloop1
        and  x0,a       a,x1
        move           a1,x:(r0)-
        move x1,a
        asr  a
```

```
diffloop1
     move x:(r0)+,a      y:(r7)+,y0
     move x:(r0)+,a      y:(r7)-,yl
     move x:(r0)-,b      a,y:(r7)+
     move                b,y:(r7)-
     eor  y0,a      a,x0
     eor  yl,b      a,xl
     and  y0,a      b,yl
     eor  yl,a      xl,b
     move           b,x:(r0)+
     move           a,x:(r0)-
     clr  a
     clr  b
     do   #4,diff2
     move           x:(r0)+,b0
     addl b,a
diff2
     rts
```

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